



# Laboratory of Nomadic Communication

Quick introduction to IEEE 802.11



# Wireless LAN Standard

A quick introduction to the IEEE 802.11 standard



# IEEE 802.11 standard

## ❑ Definition of wireless interface

- between a client and a base station (aka: Access Point, AP)
- between wireless clients (simply: stations)

## ❑ Two lower layers of the stack

- 1-PHY - radio transmission: modulations, bands, frequency, energy
- 2-MAC - medium access control: timings, retransmissions, signaling

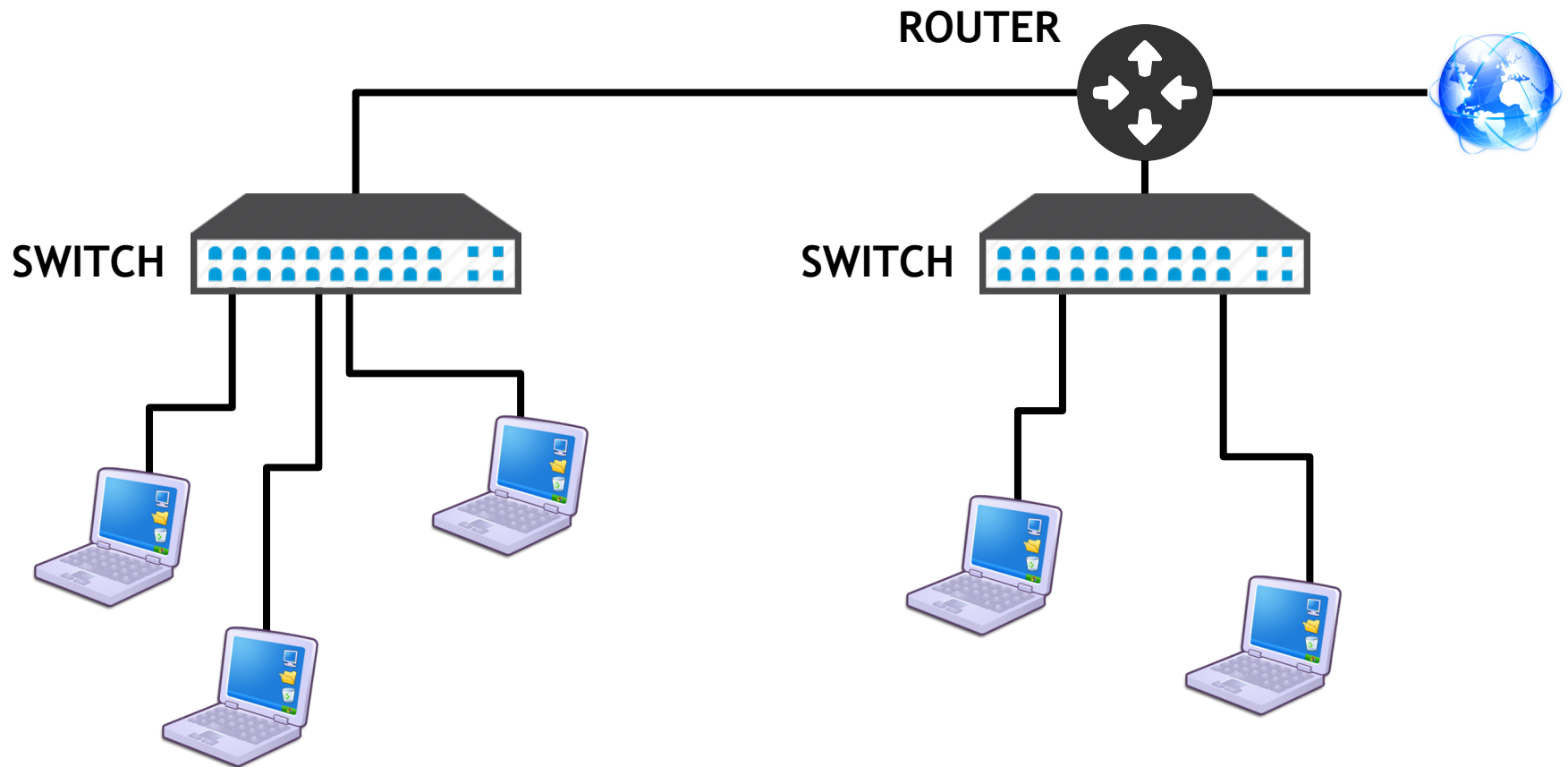
## ❑ Many amendments published since '97

- Throughput improvements (e.g., 802.11ac up to multi Gb/s)
- Security (802.11i), QoS (802.11e), reliable multicast (802.11aa)

## ❑ Very long standard: 2012 rel. (pre 11ac) is ~2800 pages!

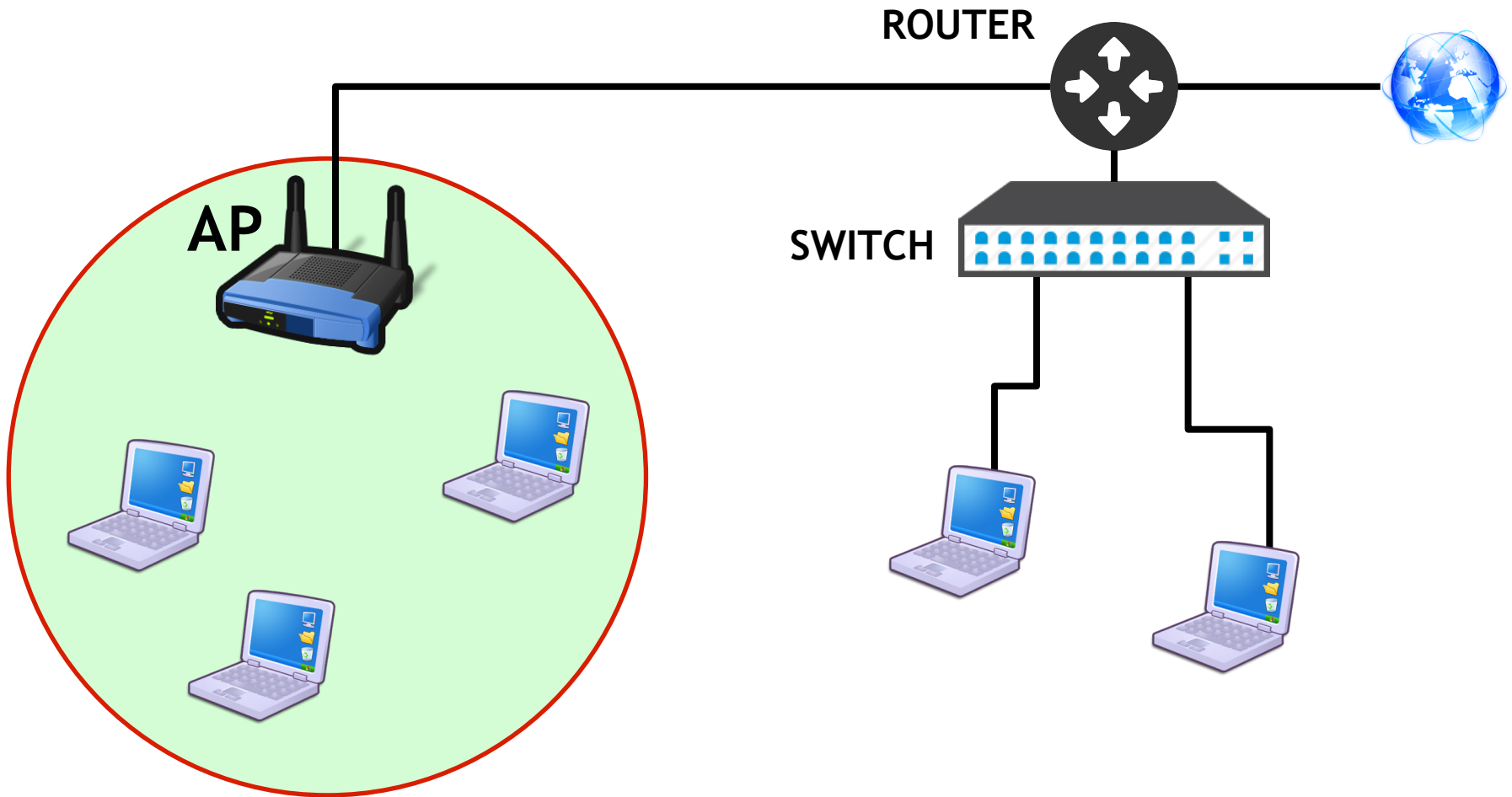


# A LAN goes WIRELESS





# A LAN goes WIRELESS/2

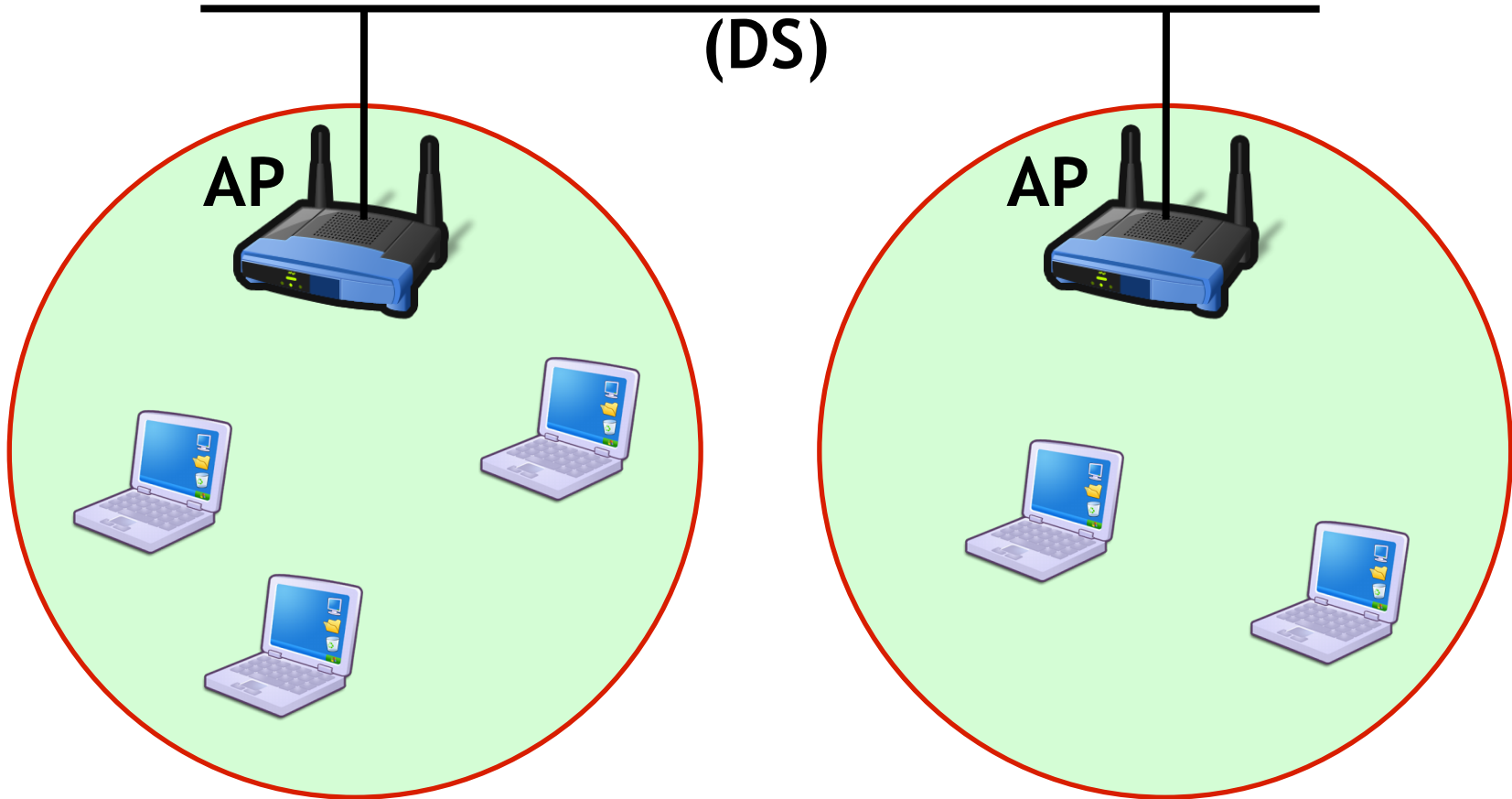




# A LAN goes WIRELESS / 3

## Distribution System

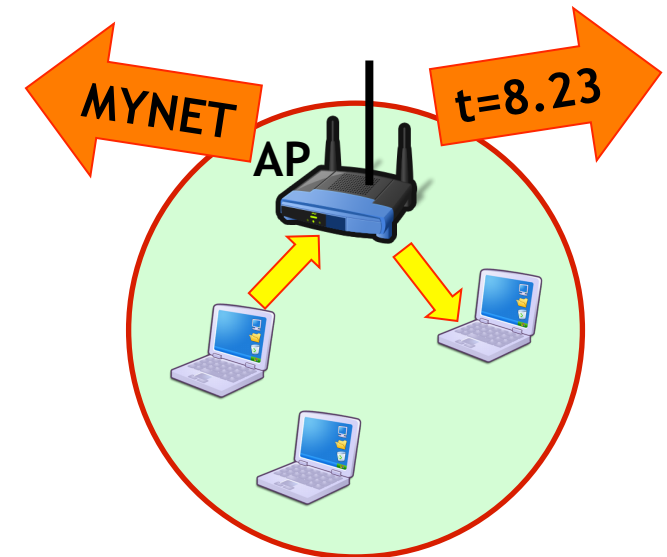
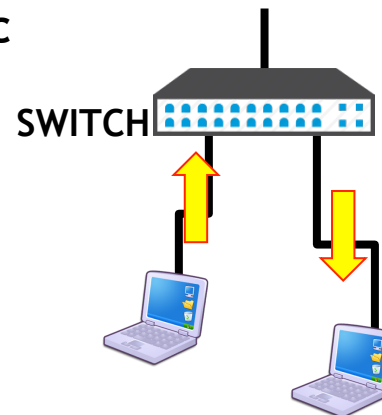
(DS)





# Wireless-LAN vs Wired-Lan

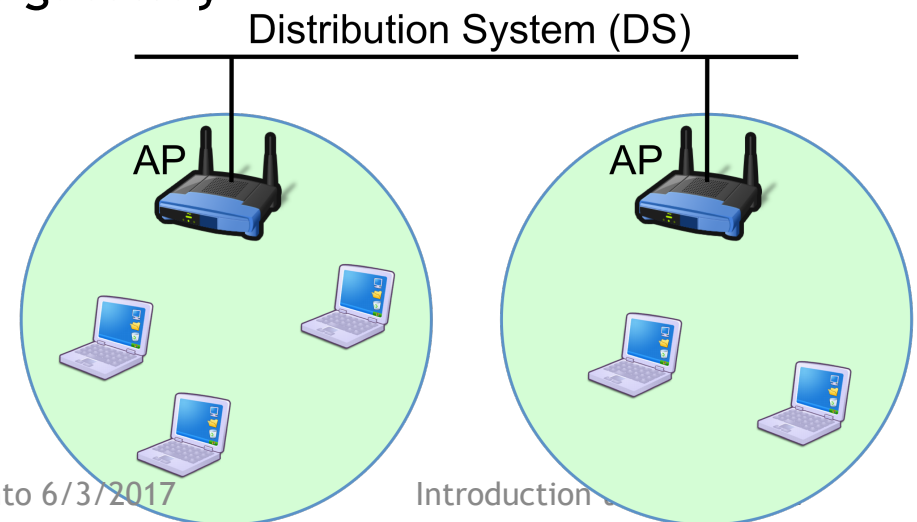
- ❑ Apparently: AP replaces the switch, air replaces cables
- ❑ AP, in fact,
  - forwards inter-stations frames (no direct communication)
  - rules stations access to the network (e.g, by authenticating)
  - manages even more issues than the switch has to, i.e.,
    - advertizes the network
    - synchronizes time etc





# 802.11 Wireless-Lan: Infrastructure Mode Basic

- ❑ Each cell is an Infrastructure “Basic Service Set” (BSS)
  - The Access Point (AP) “creates” and maintain the BSS
  - Time sync, BSS name and capabilities inside “Beacons” frame
  - All BSS traffic goes through the AP
- ❑ More cells build an “Extended Service Set” (ESS)
  - A Distribution System (DS, wired or wireless) connects all APs
  - DS may connect to an Internet gateway







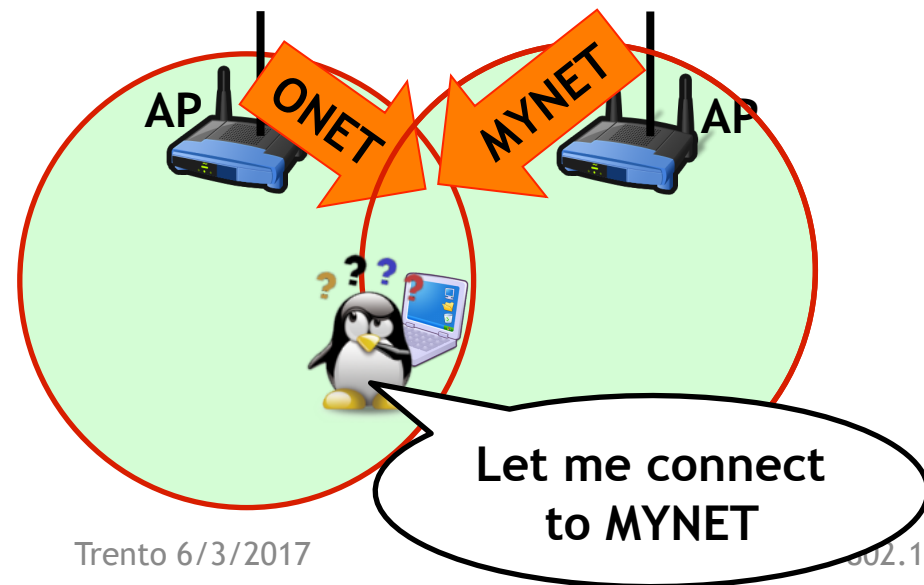
# 802.11 Wireless-Lan: Infrastructure Mode Basic

## □ A station that wants to connect does the following:

- Scan: check all the available networks for one known
  - Passively, by receiving “Beacons”, or Actively, by sending “Probes”
- Authenticate: proves to the AP she knows something
  - Easiest case: simply send her identity, wait for an ack
- Associate: station and AP shares mutual capabilities

## □ Eventually:

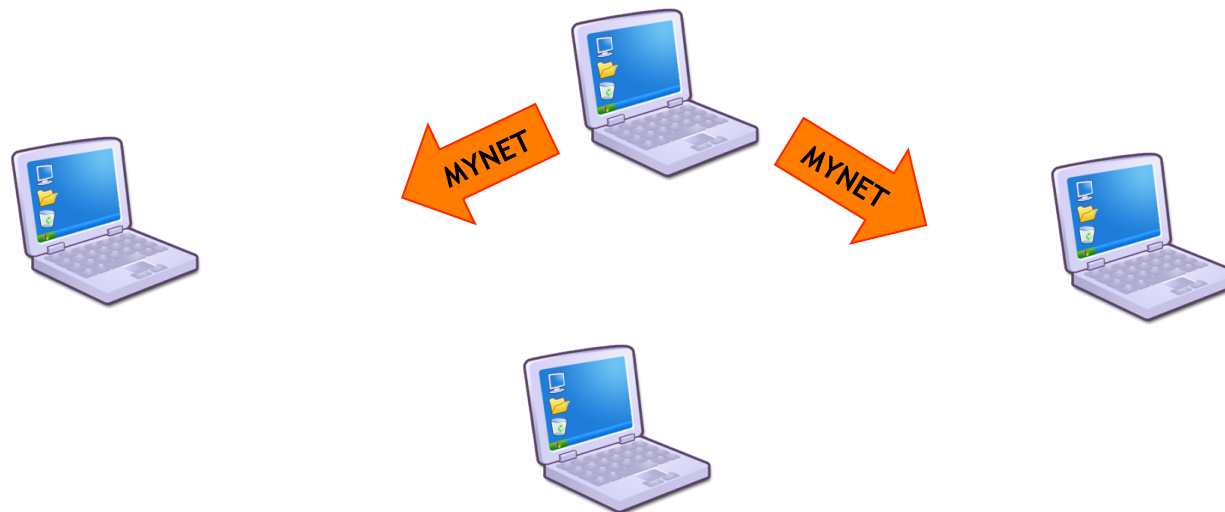
- Station sends/receives traffic





## 802.11 Wireless-Lan: Ad-Hoc Mode Basic

- ❑ Stations create an Ad-Hoc network with a specific “NAME”
  - No AP needed: Independent BSS (IBSS)
  - The first station active starts sending beacons (e.g., leader)
  - Other stations can join the Ad-Hoc network
    - If they do not receive the beacon from the leader, they transmit one!





# A very incomplete standard synopsis

Document	year	modulation add-on	band (GHz)	width (MHz)	spatial stream	Rates addon
802.11	97	DSSS	2.4	20	/	1, 2
802.11b	99	CCK	2.4	20	/	5.5, 11
802.11a	01	OFDM	5	20	/	6 ... 54 (8 rates)
802.11g	03	/	2.4	20	/	all the above
802.11n	09	MIMO & OFDM+	2.4 & 5	20, 40	up to 4	HT-PHY MCS [max 600Mb/s]
802.11ac	13	MU-MIMO & OFDM++	2.4 & 5	20, 40, 80, 160	up to 8	VHT-PHY MCS [max 6.7Gb/s]



# A very incomplete standard synopsis

□ In this course we will use 802.11bg only devices

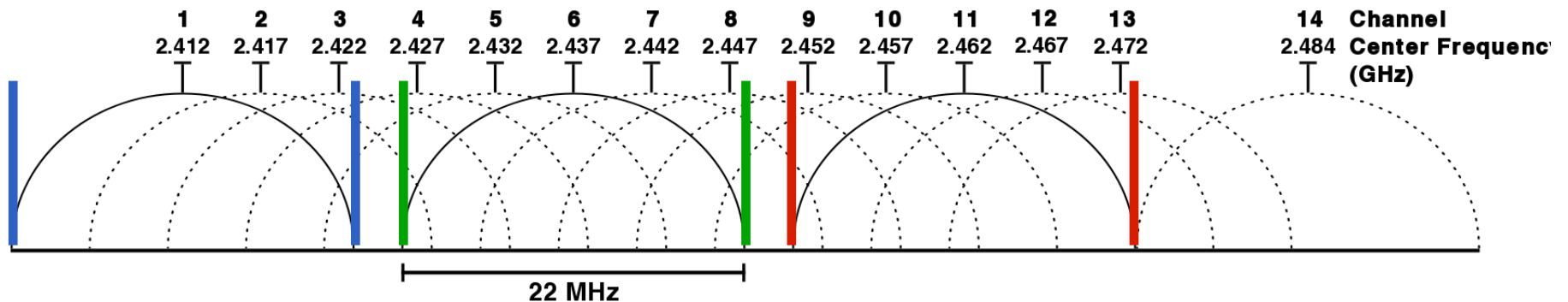
- Yeah, pretty old hardware but...
- ... we have access to the NIC internals, we can play with the standard
- All the standards share the same basic algorithms for channel access

Document	modulations	band (GHz)	width (MHz)	Rates
802.11b/g	DSSS/CCK OFDM	2.4	20	1, 2 / 5.5, 11 6, 9, 12, 18, 24, 36, 48 ,54



# IEEE 802.11: insight of the 2.4GHz band

- ❑ ISM 2.4GHz band spans range [2400-2483.5]MHz worldwide
  - Availability subject to country regulations
    - E.g., USA [1-11], Italy [1-13], Japan [1-14] (but 14 only 11b!)
- ❑ Standard: 13 channels (5 MHz spacing) + channel 14
  - $ch_N$  @ [ 2407 + 5 \* N ]MHz,  $1 \leq N \leq 13$ ; very busy ☹
  - $ch_{14}$  @ 2484MHz; not used outside Japan ☺
- ❑ How many orthogonal channels? Remember 20MHz width





# Wireless LAN Standard

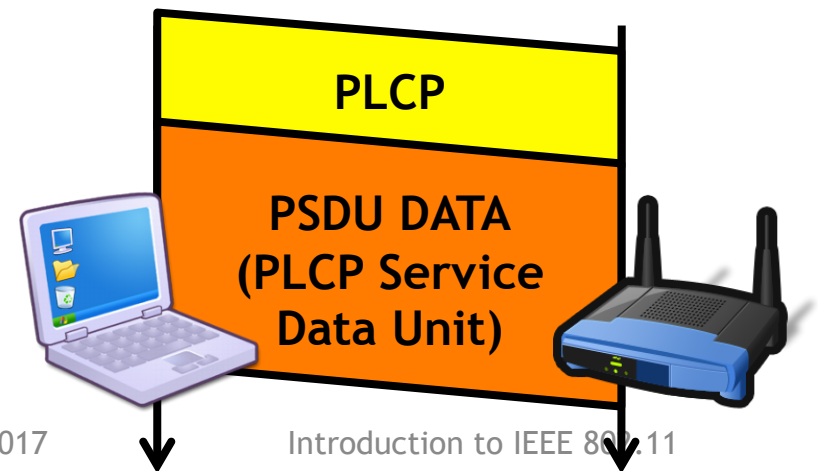
802.11bg Physical Layer analysis



# IEEE 802.11bg: a frame unit (no MAC yet)

- ❑ Each frame preceded by PLCP preamble
  - Physical Layer Convergence Procedure
- ❑ PLCP helps the receiver
  - Understanding a transmission is beginning (energy raise)
  - Knowing which data-rate encoding is used for data and its length
  - Synchronizing the decoding subsystem
- ❑ PLCPs of 11b and 11g differ
  - Let's check!

PPDU  
(PLCP Protocol  
Data Unit)

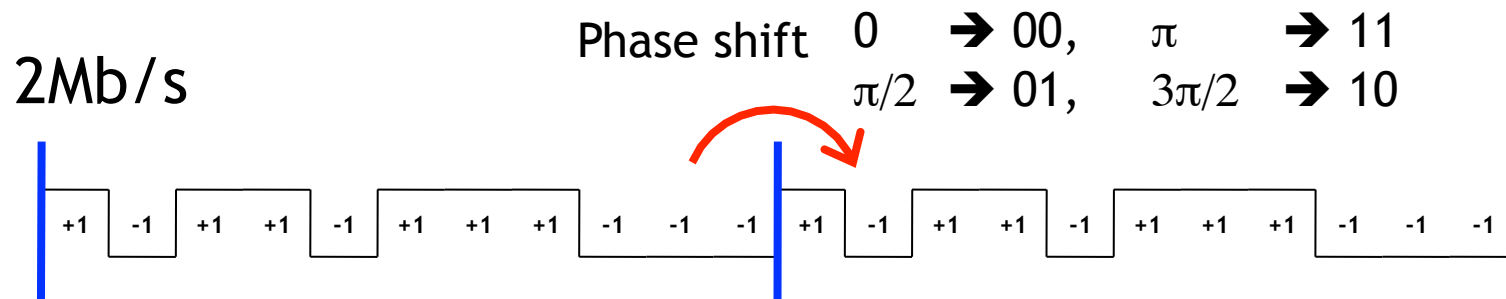
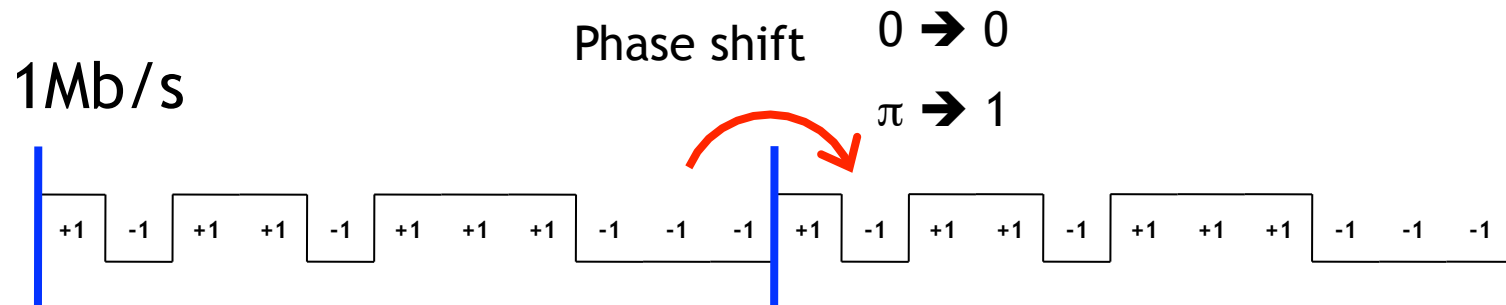




# IEEE 802.11b: PPDU format/1

## □ Direct Sequence Spread Spectrum modulations

- Each bit is “spread” to a sequence of 11 chips (Barker train)
- A sequence of 11 “chips” transmitted repeatedly by shifting phase
- Phase shift of consecutive chip trains encode the PPDU



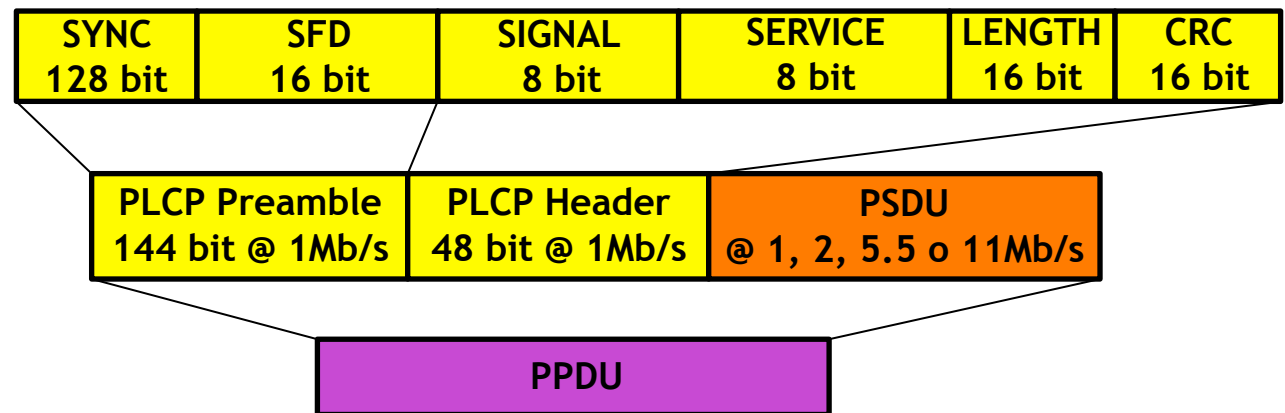




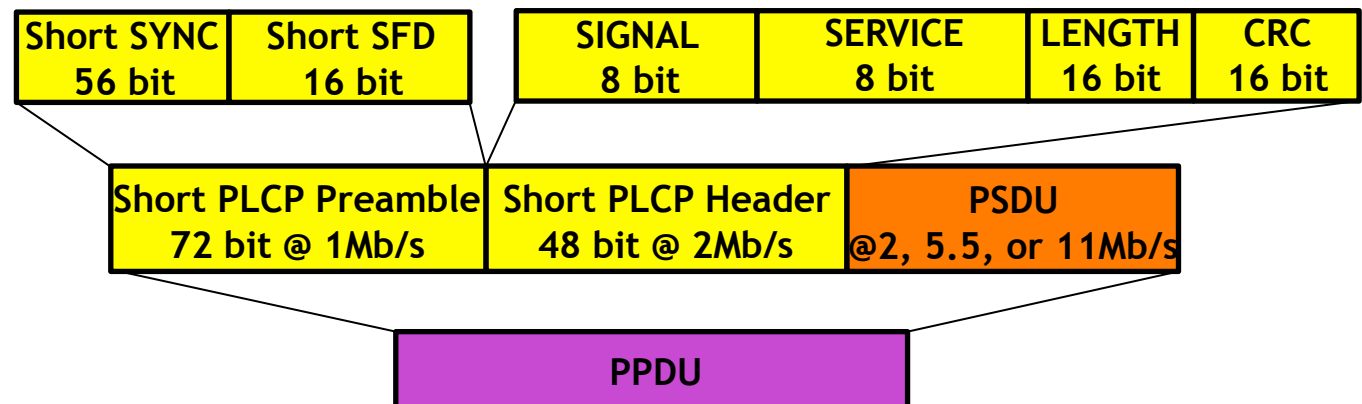
# IEEE 802.11b: PPDU format/2

□ Two possible PLCP format:

Long PLCP is 192 $\mu$ s



Short PLCP is 96 $\mu$ s  
For 2, 5.5, 11Mb/s





## IEEE 802.11b: PPDU format/3

❑ PPDU format is very inefficient!

❑ E.g.: Acknowledgement (shortest frame), 14byte = 112bit

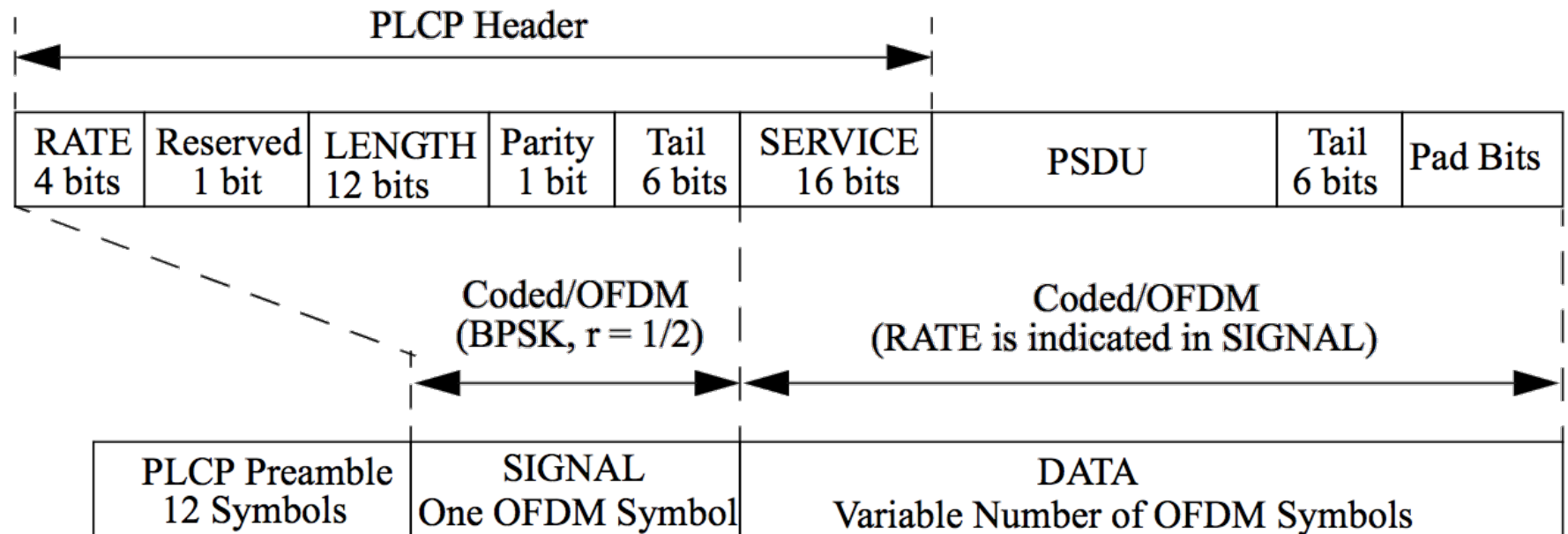
- |             |                     |             |
|-------------|---------------------|-------------|
| - Preamble  | PSDU                |             |
| - @1Mb/s,   | 192 $\mu$ s         | 112 $\mu$ s |
| - @2Mb/s,   | 192 $\mu$ s (or 96) | 56 $\mu$ s  |
| - @5.5Mb/s, | 192 $\mu$ s (or 96) | 21 $\mu$ s  |
| - @11Mb/s,  | 192 $\mu$ s (or 96) | 11 $\mu$ s  |

**PLCP much longer than actual data!**



# IEEE 802.11g: PPDU format/1

- 802.11g: new PLCP, very short



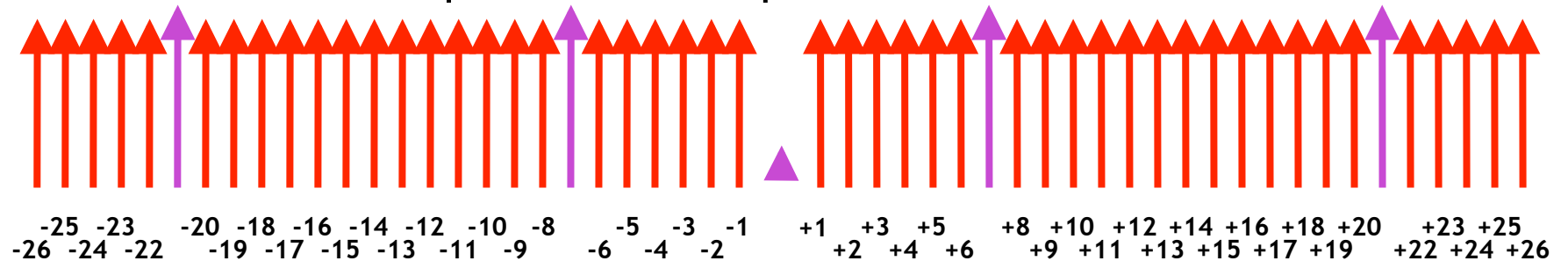
PLCP is 20 $\mu$ s



# IEEE 802.11g: PPDU format/2

## ❑ Orthogonal Frequency Division Multiplexing (6Mb/s)

- PPDU is divided in groups of 24 bits (three bytes)
- Each group of 24 is expanded to 48 bit (FEC)
- Each bit of these 48 weights a specific carrier (with -1 for 0, +1 for 1)
- Five additional pilot carriers - pink - inserted



- Time signal computed by running ifft of the carrier weights
- Each group takes  $4\mu\text{s}$ , call this “OFDM symbol”

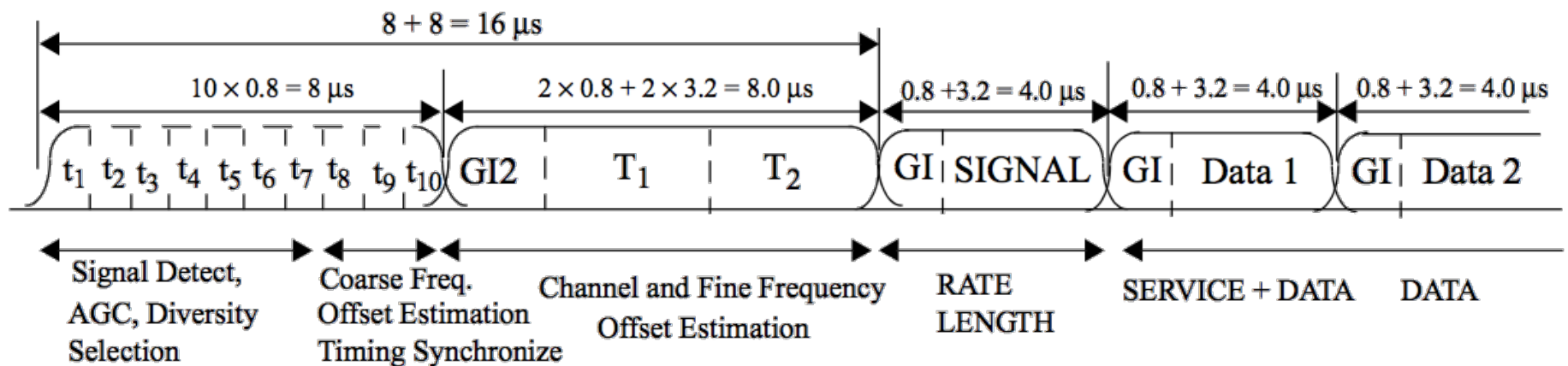
## ❑ Higher data rates map more bits to each carrier/symbol



# IEEE 802.11g: PPDU format/2

## □ 802.11g: frame includes

- PLCP Preamble made of 10 short symbols ( $8\mu s$ ) + 2 long ones ( $8\mu s$ )
- PLCP Header made of
  - SIGNAL field in its own symbol ( $4\mu s$ , PSDU length+rate)
  - SERVICE field, first 16 bits of first data symbol
- PSDU made of symbols, each one carrying N bits, N depends on Rate
- PSDU terminate with CRC32 and at least 6 padding bits





# IEEE 802.11g: PPDU format/2

## □ 802.11g: data payload

- Bit expanded by convolutional encoder for FEC,  $R = [1/2, 2/3, 3/4]$
- Groups of  $N_{CBPS}$  (Coded Bit Per Symbol) or  $N_{DBPS}$  (Data Bit Per Symbol)
- Each subcarrier transport  $N_{BPSC}$  bit (Bit Per SubCarrier)

Modulation	R	$N_{BPSC}$	$N_{CBPS}$	$N_{DBPS}$	Data rate
BPSK	1/2	1	48	24	6
BPSK	3/4	1	48	36	9
QPSK	1/2	2	96	48	12
QPSK	3/4	2	96	72	18
16-QAM	1/2	4	192	96	24
16-QAM	3/4	4	192	144	36
64-QAM	2/3	6	288	192	48
64-QAM	3/4	6	288	216	54



## IEEE 802.11g: frame format/6

- E.g.: Acknowledgement, 14byte = 112bit
- PLCP:  $20\mu s$
- $DATA_{PSDU}: 16b(\text{SERVICE})+112b(\text{PSDU})+6b(\text{tail}_{\min})=134b$

Data rate	PLCP	$N_{DBPS}$	bit	symbol	$\Delta T$	Extension	Total
6	$20\mu s$	24	134	6	$24\mu s$	$6\mu s$	$50\mu s$
9	$20\mu s$	36	134	4	$16\mu s$	$6\mu s$	$42\mu s$
12	$20\mu s$	48	134	3	$12\mu s$	$6\mu s$	$38\mu s$
18	$20\mu s$	72	134	2	$8\mu s$	$6\mu s$	$34\mu s$
24	$20\mu s$	96	134	2	$8\mu s$	$6\mu s$	$34\mu s$
36	$20\mu s$	144	134	1	$4\mu s$	$6\mu s$	$30\mu s$
48	$20\mu s$	192	134	1	$4\mu s$	$6\mu s$	$30\mu s$
54	$20\mu s$	216	134	1	$4\mu s$	$6\mu s$	$30\mu s$



# Wireless LAN Standard

802.11bg Frame format analysis (@ layer 2)

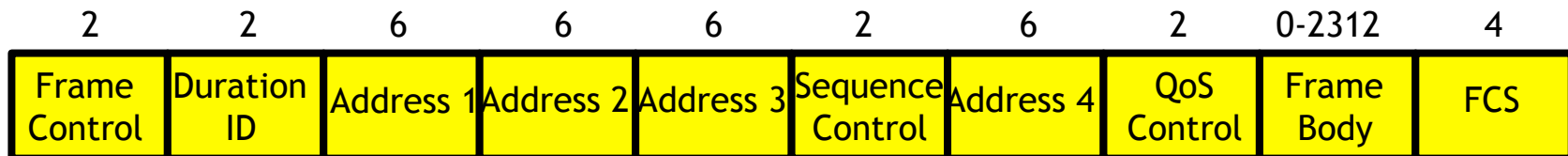




# IEEE 802.11: Frame types

## □ Three types of Mac Protocol Data Unit (MPDU):

- Management, e.g. Association Request/Response, Beacon, (De)Auth
  - Network/BSS Advertisement, BSS Join, Authentication etc
- Control, e.g. ACK, RTS, CTS, Poll, etc
  - For channel access (RTS, CTS), positive frame acknowledgment
- Data: Plain data + QoS Data, etc
  - Frames with user data
- MPDU fields: depend on frame type!

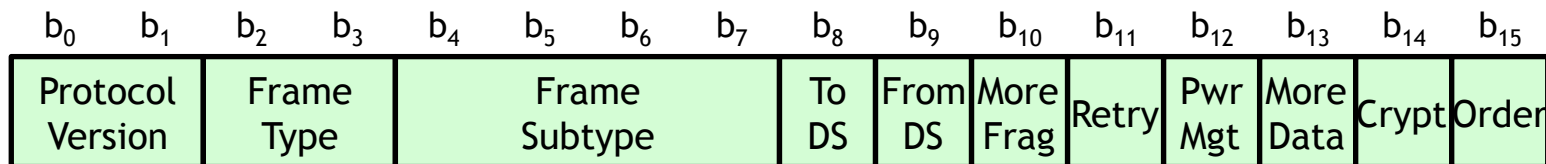




# IEEE 802.11: PSDU fields/1

## □ Frame Control:

- Protocol version, only 0 today
- Type and Subtype encode frame type + subtype
- ToDS: frame is for Distribution System; FromDS frame is from DS
  - If both set to 1, frame is transported by a Wireless DS
- More: announce other fragments are coming (PSDU is fragmented)
- Retry: help rx'er understanding this is a retransmission
- {Pwr Mgt, More Data} deal with power management, save
- Protected: announce Frame Body is encrypted





# IEEE 802.11: PSDU fields/2

## ❑ Duration/ID

- Meaning depends on MPDU type
- Data: number of  $\mu s$  after frame end during which medium is reserved
  - Used by Virtual Carrier Sense, set the Network Allocation Vector (NAV)

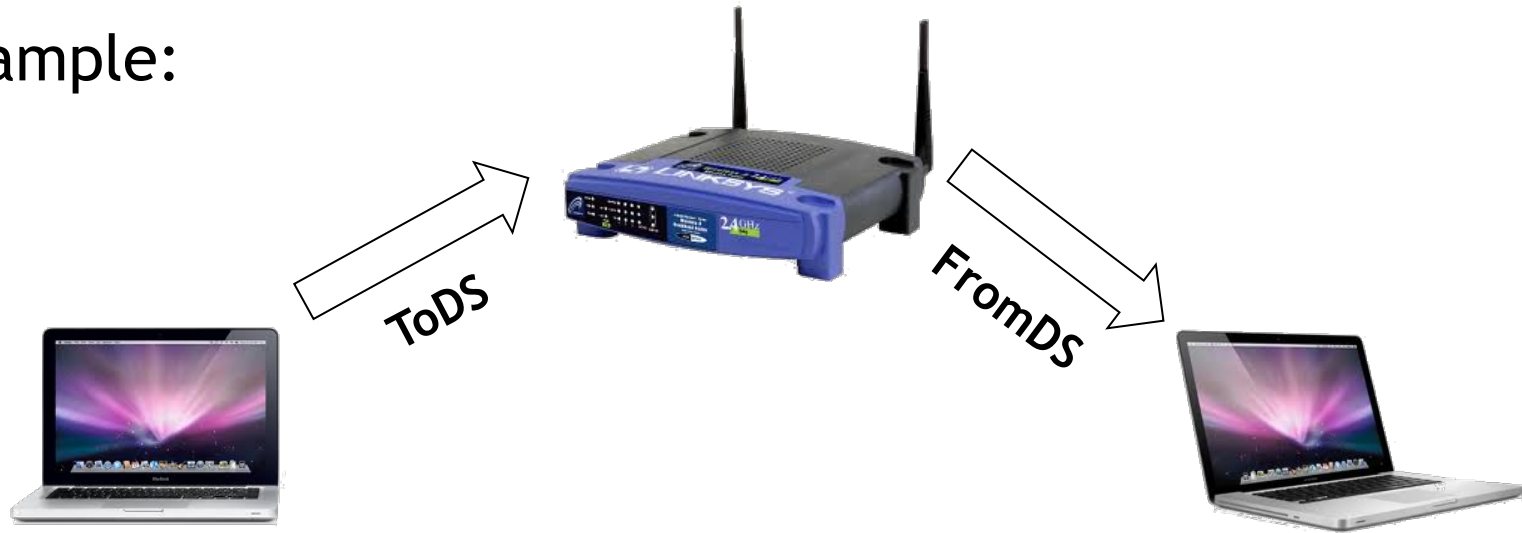
## ❑ Address fields: they depends on ToDS/FromDS fields:

- BSSID: Basic Service Set IDentification
  - Address of the AP
- DA: Destination Address, “final destination”
- RA: Receiver Address, immediate frame destination
- SA: Source Address, who has generated this frame
- TA: Transmitter Address, who has forwarded this frame



# IEEE 802.11: PSDU fields/3

□ Example:



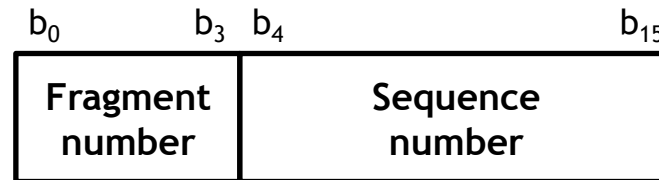
To DS	From DS	Address 1	Address 2	Address 3	Address 4
0	0	RA = DA	TA = SA	BSSID	N/A
0	1	RA = DA	TA = BSSID	SA	N/A
1	0	RA = BSSID	TA = SA	DA	N/A
1	1	RA	TA	DA	SA



## IEEE 802.11: PSDU fields/4

### ❑ Sequence Control:

- Fragment number, 4 bits
  - For fragmented PSDU, it's the number of this fragment
- Sequence Number, 12 bits, unique for PSDU
  - Identify the PSDU (used by rx'er to avoid accepting same frame > once)



### ❑ QoS Control: identify Traffic Category (optional field)

### ❑ FCS: CRC/32 Frame Check Sequence protecting the PSDU

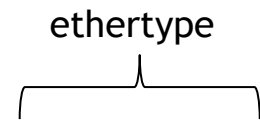


# IEEE 802.11: PSDU example - Data Frame

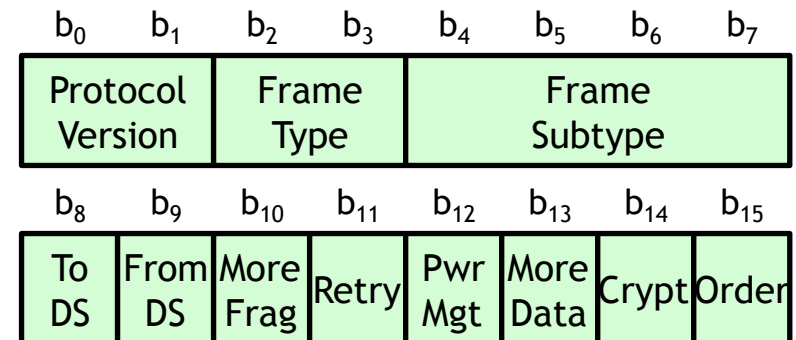
## □ IP packet, no QoS, from STA to AP (ToDS): Data frame

- Logical Link Control (LLC) encapsulation is used

- 8 bytes before IP: 0xAA, 0xAA, 0x03, 0x00, 0x00, 0x00, 0x08, 0x00



- Type: Data frame  
SubType: 0 ⇒ Byte#0 := 0x08

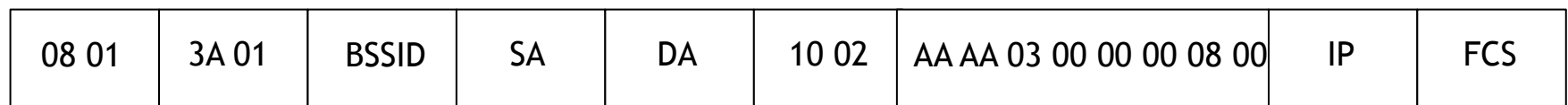
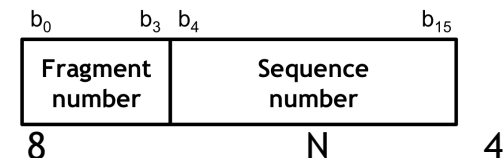


- ToDS ⇒ Byte#1 := 0x01

- Duration: time to tx an ACK + SIFS

- Address: it's a ToDS frame, fill the three address fields

- SeqCTRL: seq. no:=33 ⇒ SeqCTRL:=0x0210



30



# IEEE 802.11: PSDU example - Data Frame

## □ Example with WireShark

- Open a trace file
- Show Beacons
- Show Data, retry etc



# Wireless LAN Standard

The Basic Access Scheme

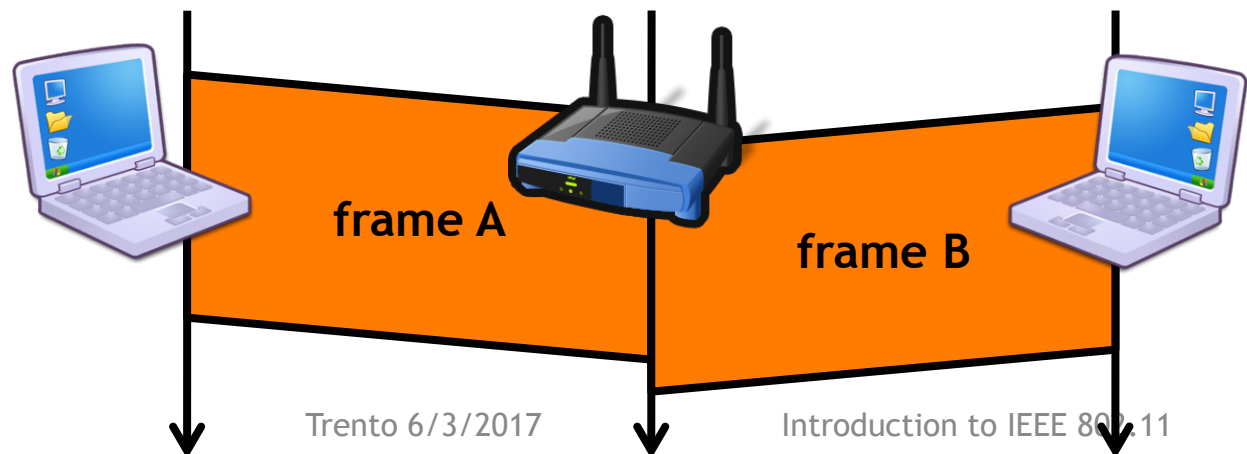
Distributed Coordination Function (DCF)





# Transmissions in a Broadcast Medium

- ❑ If two stations transmit at the same time
  - Receiver(s) cannot decode packets (collision)
  - Transmitters do not know whether data was received
- ❑ 802.11 standard introduces
  - Slotted medium
  - Positive Acknowledgment with Retransmission
  - Carrier Sense Multiple Access/Collision Avoidance





# Time Slot

- ❑ Time is divided into intervals, called **slots**
  - Working with slot (tx may start with slot) reduces uncertainty
  
- ❑ A Slot is the system unit time
  - 802.11b Slot Time is  $20\mu\text{s}$ , 11g is  $9\mu\text{s}$
  
- ❑ Time synchronized by Beacons transmitted by the BSS AP
  - Each Beacon carries a 64bit time value ( $1\mu\text{s}$  granularity)
  - Stations in the BSS copy beacon time to their clock registers
  - Skews due to poor clock design periodically corrected
  
- ❑ A BSS is a synchronous system!!



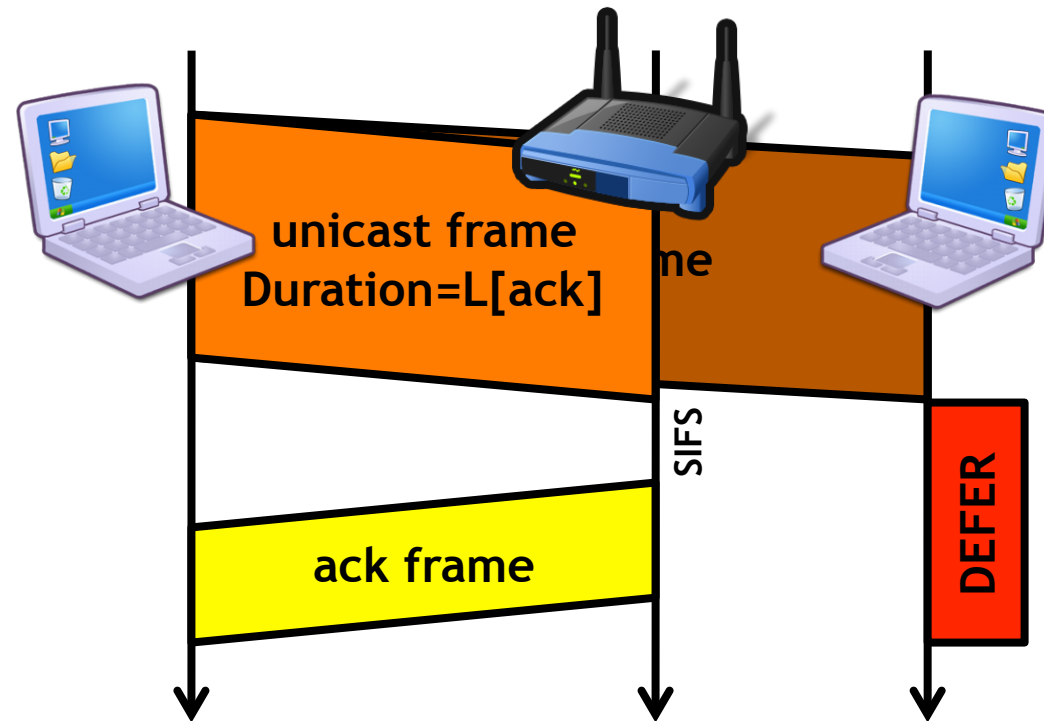
# InterFrame Space (IFS)

- ❑ Time interval between consecutive transmissions
- ❑ Different IFSs allow different access priorities
  - Short IFS: separate transmissions belonging to the same “dialogue”
    - SIFS in 11bg it's  $10\mu\text{s}$
  - Point Coordination IFS: used by the Point Coordinator
    - PIFS is  $\text{SIFS} + \text{Slot Time}$
  - Distributed IFS: waited by stations when contending for a free channel
    - DIFS is  $\text{SIFS} + 2 * \text{Slot Time}$
  - Extended IFS: waited by stations when receiving a bad frame
    - EIFS is  $\text{SIFS} + \text{TxTime}[\text{AckFrame}] + \text{DIFS}$



# Positive Acknowledgment & Retransmissions

- ❑ Received unicast frame must be acknowledged
  - Other nodes defer transmissions by using the received “Duration”





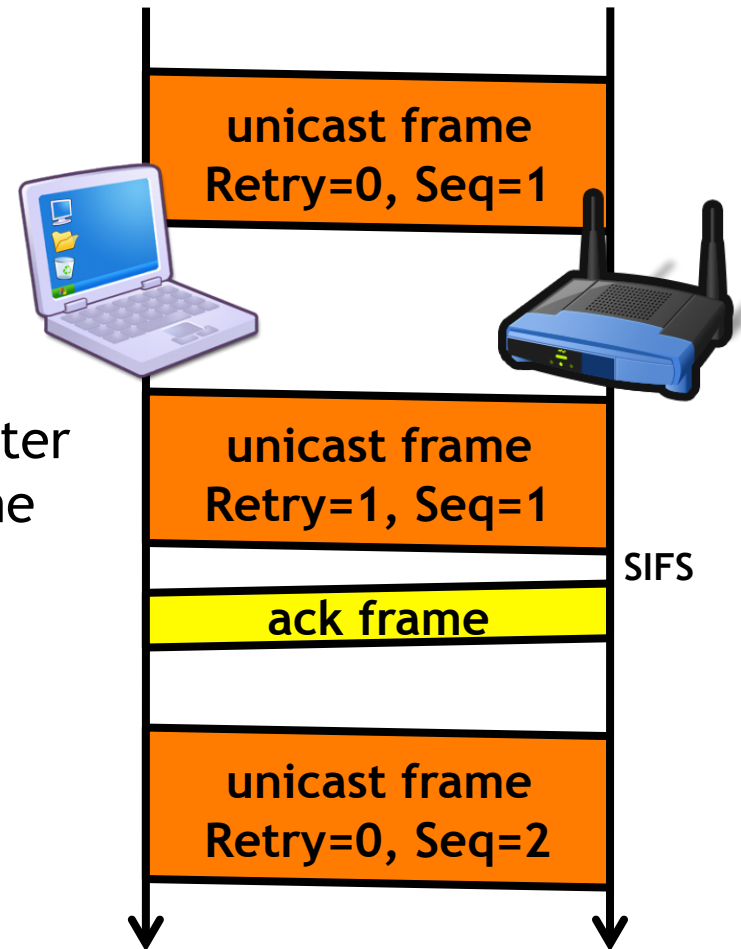
# Positive Acknowledgment & Retransmissions/2

## ❑ If no acknowledgment coming from receiver

- Retransmit the frame with Retry bit set and same sequence counter

## ❑ When acknowledgment received

- Increase the sequence counter and switch to the next frame





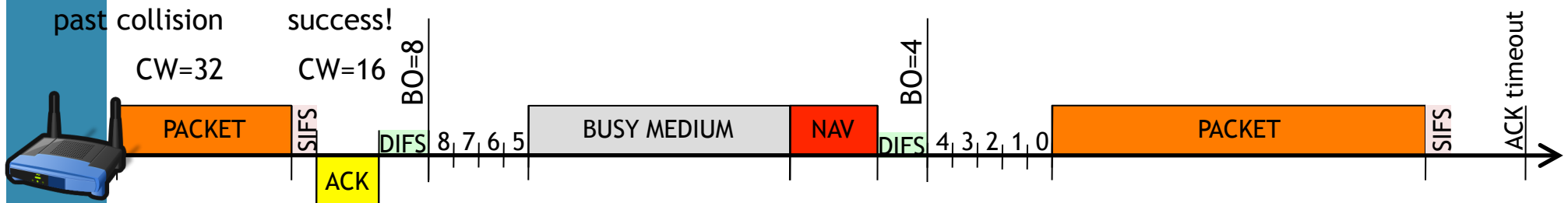
# Carrier Sense Multiple Access/ Collision Avoidance

- ❑ Apart from slot synchronization
  - No other explicit coordination among stations
- ❑ To **avoid** repeated collisions
  - Carrier Sense Multiple Access/Collision Avoidance (CSMA/CA)
  - Given free-space attenuation, /Collision Detection is not feasible!
- ❑ CSMA/CA Basic
  - Stations willing to transmit have to contend for channel access
  - A station repeats the contention procedure for every (re)transmission



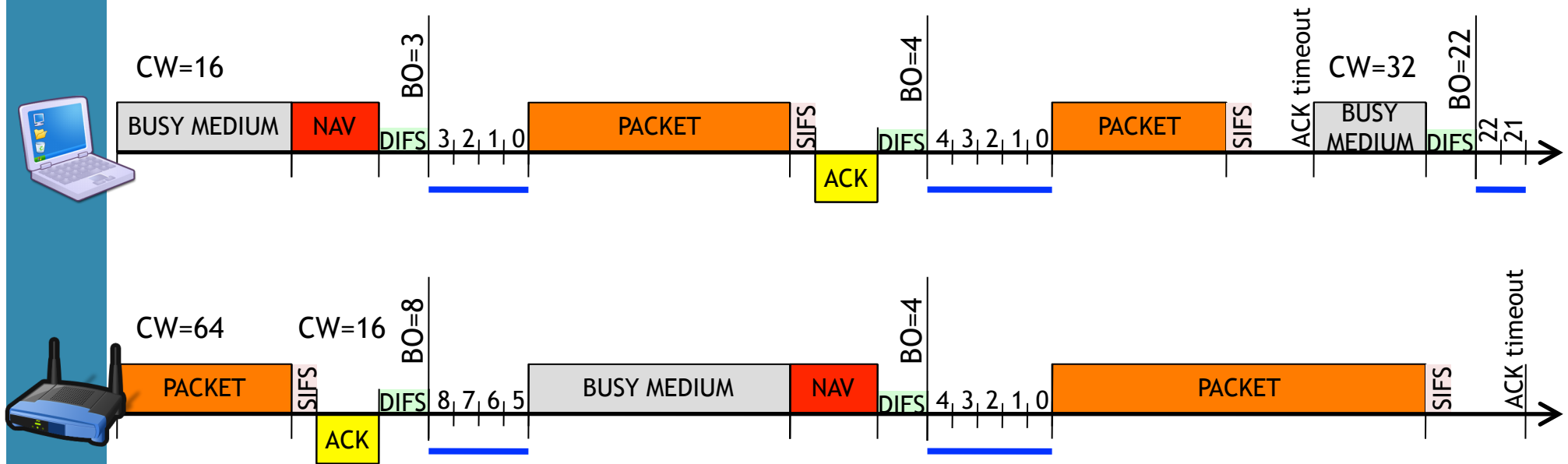
# CSMA/CA and channel contention

- ❑ Each station keeps a Contention Window (CW) parameter
  1. At the end of the previous transmission attempt
    - If collision (no ack), double CW, otherwise reset to  $CW_{min}$
    - Extract Backoff value (BO)  $\in U[0, CW - 1]$
  2. “Monitor channel free for  $t > DIFS$ ”
  3. Backoff stage: decrement BO to zero
    - Backoff: if medium free, decrement BO at every SLOT
    - When medium busy  $\Rightarrow$  Suspend: BO freezed & goto 2
    - If BCKOFF == 0  $\Rightarrow$  Transmit!





# CSMA/CA and channel contention/2

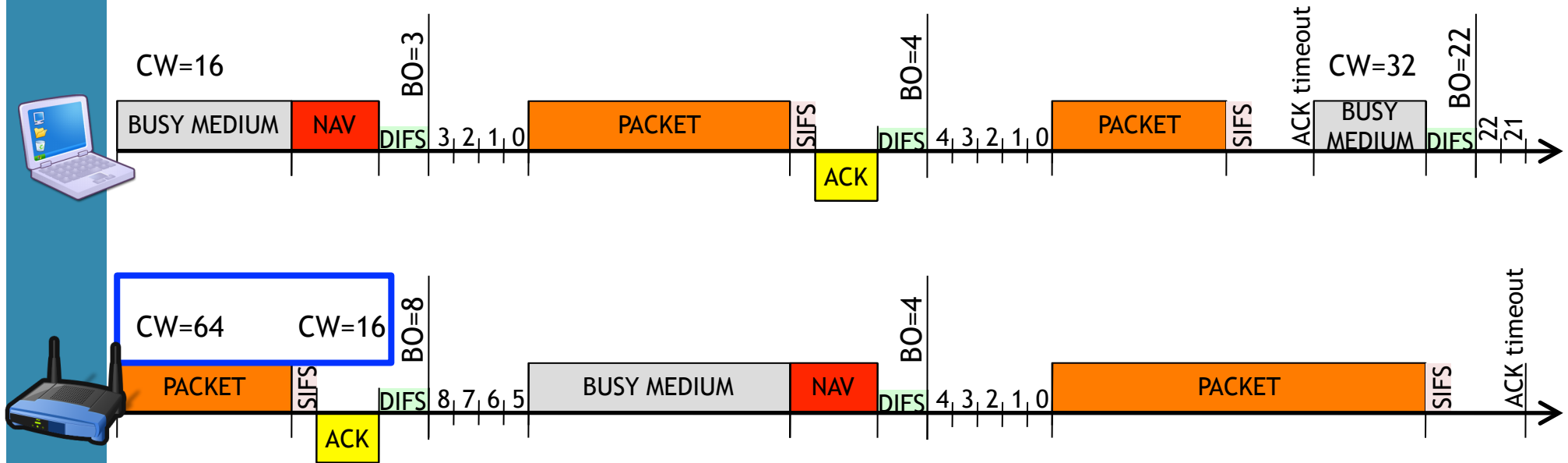


**Backoff count-down,  
channel idle**





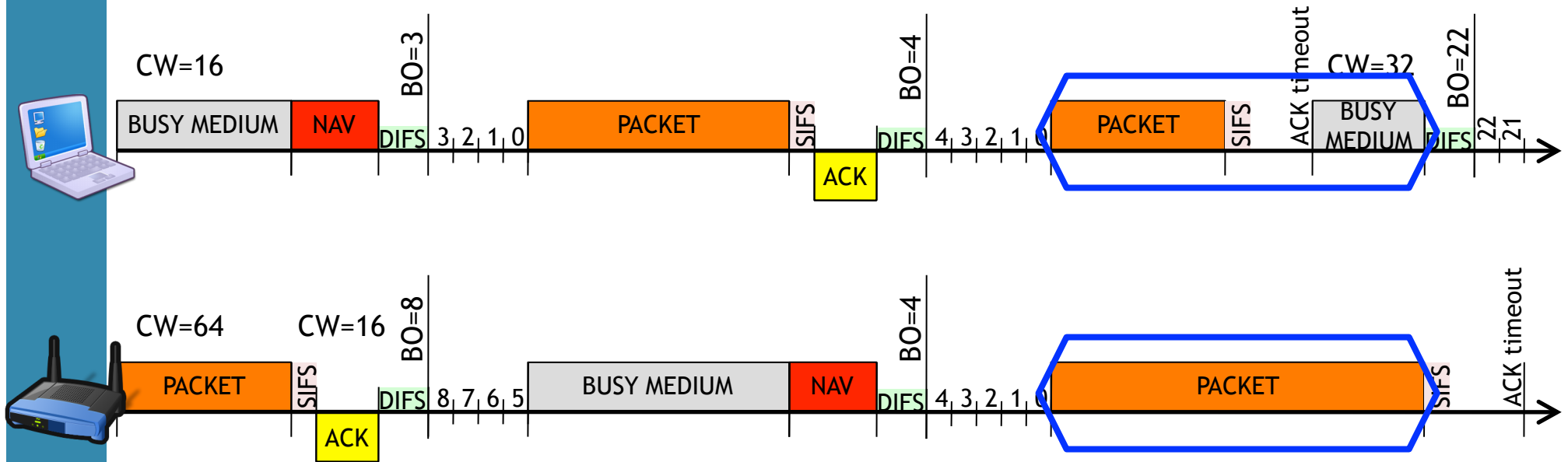
# CSMA/CA and channel contention/3



**CW reset after success**



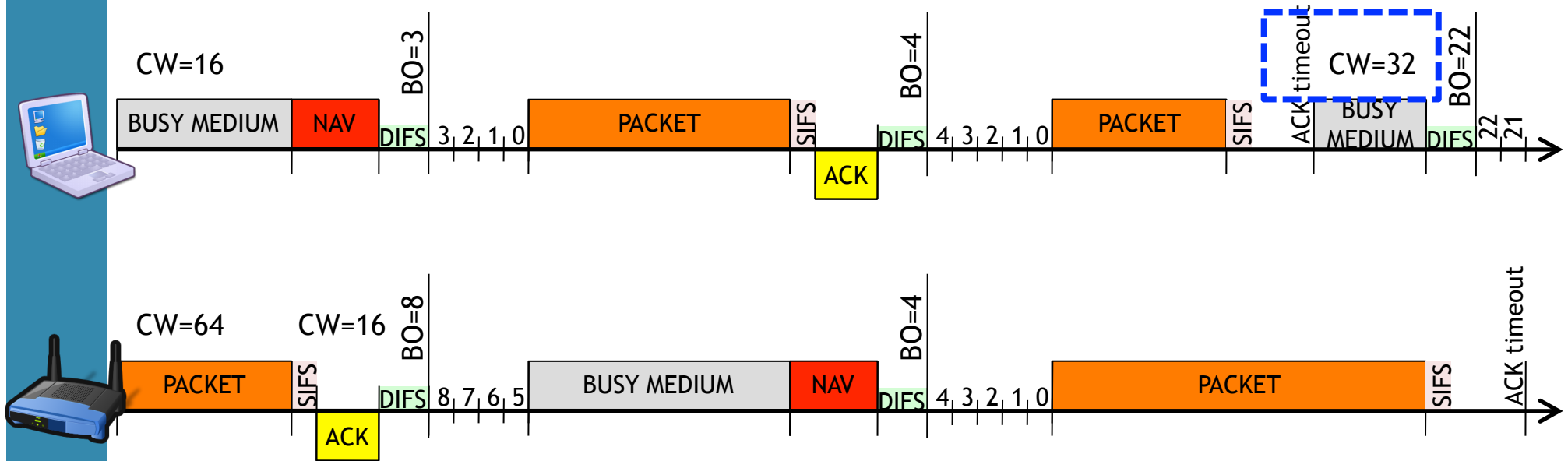
# CSMA/CA and channel contention/4



**Collision!**



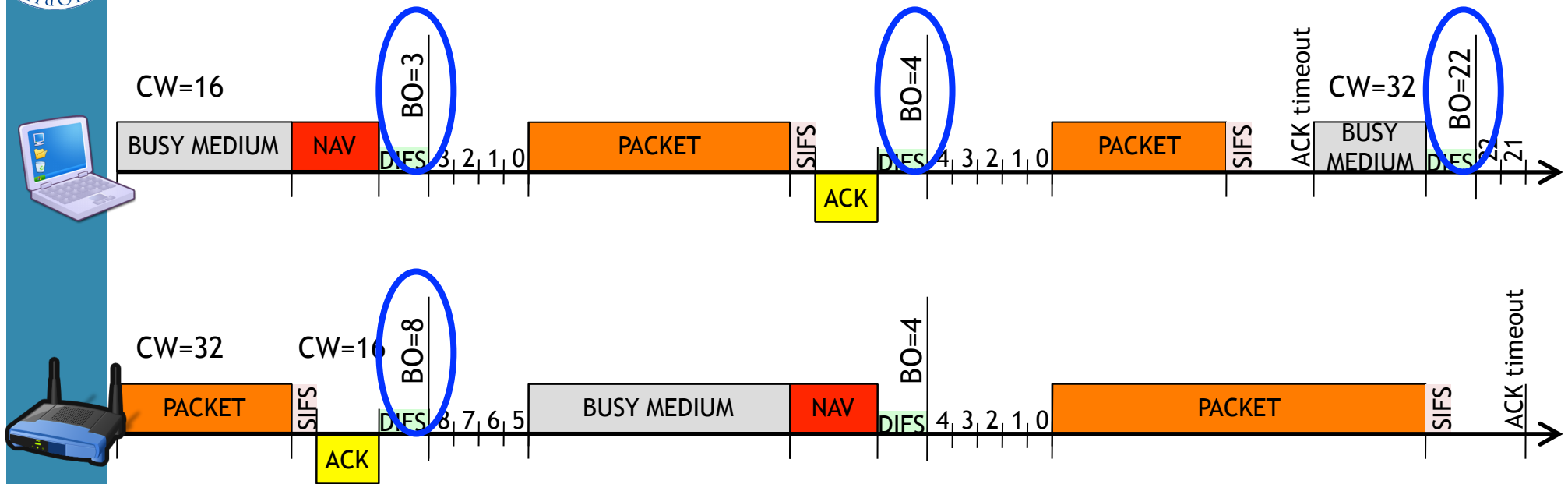
# CSMA/CA and channel contention/5



CW doubling after failure



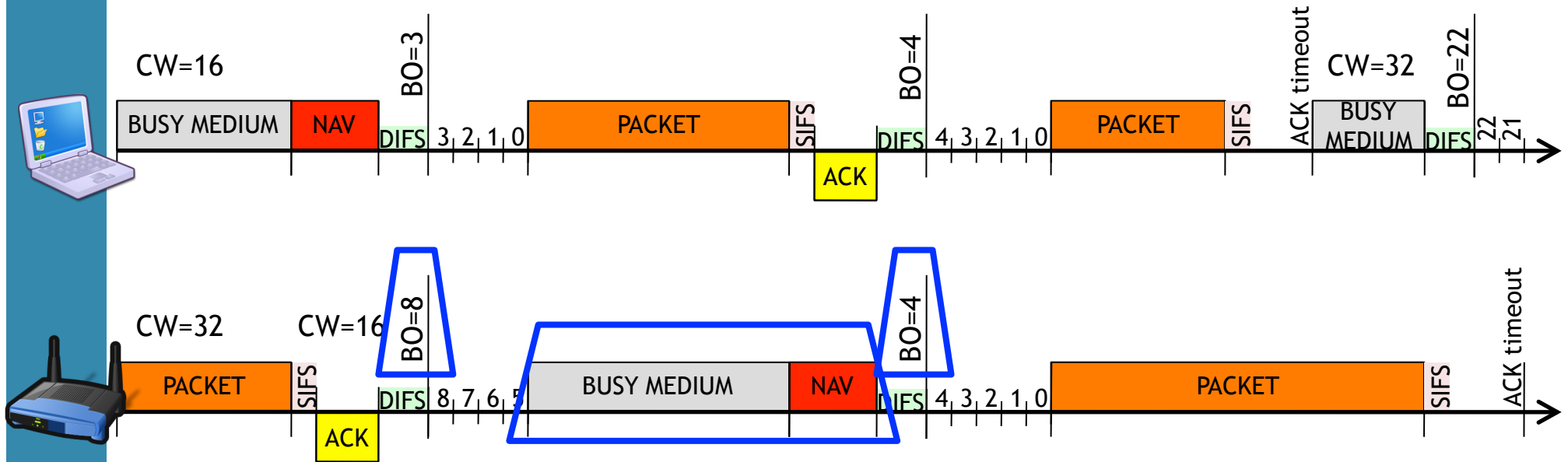
# CSMA/CA and channel contention/6



**Backoff Extraction**



# CSMA/CA and channel contention/7



**Backoff Freeze,  
channel busy**



## CSMA/CA, Exponential Backoff rule

- ❑ BCKOFF value is computed after every tx attempt
  - BCKOFF taken from  $[0, 1, \dots, CW-1]$  with uniform distribution
- ❑ Contention Window (CW) refreshed
  - $CW = 2 * CW$  if after tx attempt there is a collision
    - Up to  $CW_{max}$ , then stay with  $CW_{max}$
  - $CW = CW_{min}$  if after tx attempt, tx was acked by acknowledgment
- ❑ Standard values:
  - $CW_{min} = 16/32$ ,  $CW_{max} = 1024$
- ❑ For tx a packet that requires ACK
  - Repeat access procedure up to  $MAX_{times}$  (e.g., 7), then discard packet
- ❑ This procedure guarantees network works correctly!!



# CSMA/CA, pseudo-code

## ❑ Neglecting initializations:

procedura di trasmissione

```
while true do
  wait for packet
  wait for channel_idle_for_DIFS
  while channel_is_idle do ←
    BCKOFF--
    if BCKOFF == 0 then
      send packet
      wait for acknowledgement or timeout
      if ack received then
        reset CW
        remove packet
      else
        grow CW
      end if
      BCKOFF = extract CW
    leave do
  end fi
end do
end do
```



# Wireless LAN Standard

Rate control algorithm (super-quick)





## IEEE 802.11bg: rate choice

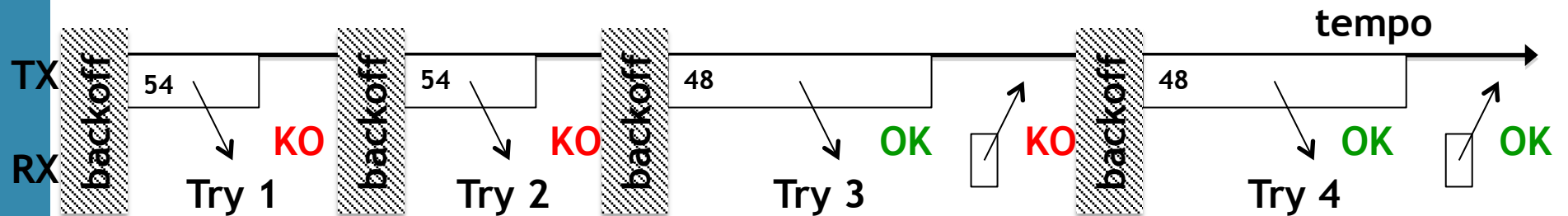
- ❑ How to choose the rate is not specified by the standard
  - Rate Controller algorithm: RC
- ❑ RCs use feedback based techniques
- ❑ E.g. Minstrel algorithm, the default today in Linux kernel
  - Count total frames transmitted PER every rate, assess success probability
  - Rate that has best success delivery ratio is the winner
  - Periodically (every N frames) send a frame at a “look-around” rate
    - Constantly scan the entire rate set
  - Rely on frames that require ACK, by counting:
    - Number of attempts per packet
    - Failed rate, success rate



# IEEE 802.11: rate choice/2

## Example: UDP packet

- RC set up these rates:  $[54\text{Mb/s}^{\{1,2\}}, 48\text{Mb/s}^{\{3,4\}}, 12\text{Mb/s}^{\{5\}}, 1\text{Mb/s}^{\{6,7\}}]$



- At the end of this packet, RC refreshes its table...

Rate	Success	Failure	Rate	Success	Failure
54	2812/3004 (93%)	192/3004 (7%)	54	2812/3006 (93%)	194/3006 (7%)
48	408/507 (80%)	99/507 (20%)	48	409/509 (80%)	100/509 (20%)
36	102/402 (25%)	300/402 (75%)	36	102/402 (25%)	300/402 (75%)

- Don't change decision (not now 😊 )



## Bibliography

- ❑ [1] IEEE 802.11-2007, Wireless LAN Medium Access Control (MAC) and Physical Layer (PHY) Specifications, June 2007.
- ❑ [2] Tutorial on 802.11n from Cisco:  
[http://www.wireshark.ch/download/Cisco\\_PSE\\_Day\\_2009.pdf](http://www.wireshark.ch/download/Cisco_PSE_Day_2009.pdf)
- ❑ [3] G. Bianchi, “Performance analysis of the IEEE 802.11 distributed coordination function”. IEEE Journal on Selected Areas in Communications, 18(3), pp. 535-547, 2000.