



STAGE WEEK 2015

Playing with Wireless Networking

Medium Access Control algorithms

Introduction to IEEE 802.11





Wireless LAN Standard

A quick introduction to the IEEE 802.11 standard

IEEE 802.11 standard

❑ Definition of wireless interface

- between a client and a base station (aka: Access Point, AP)
- between wireless clients (simply: stations)

❑ Two lower layers of the stack

- 1-PHY - radio transmission: modulations, bands, frequency, energy
- 2-MAC - medium access control: timings, retransmissions, signaling

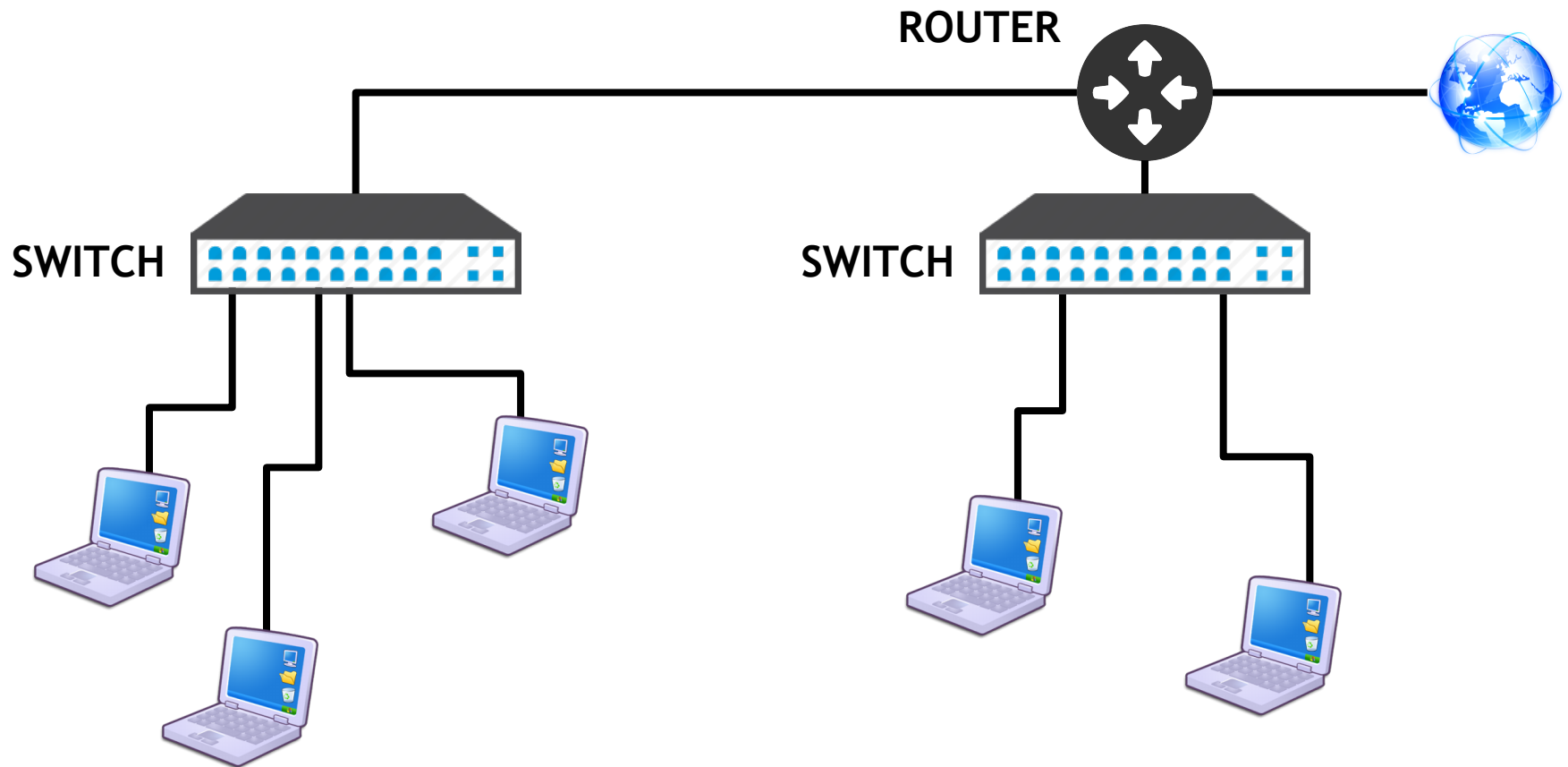
❑ Regulator published many amendments since '97

- Throughput improvements (e.g., 802.11ac up to multi Gb/s)
- Security (802.11i), QoS (802.11e), reliable multicast (802.11aa)

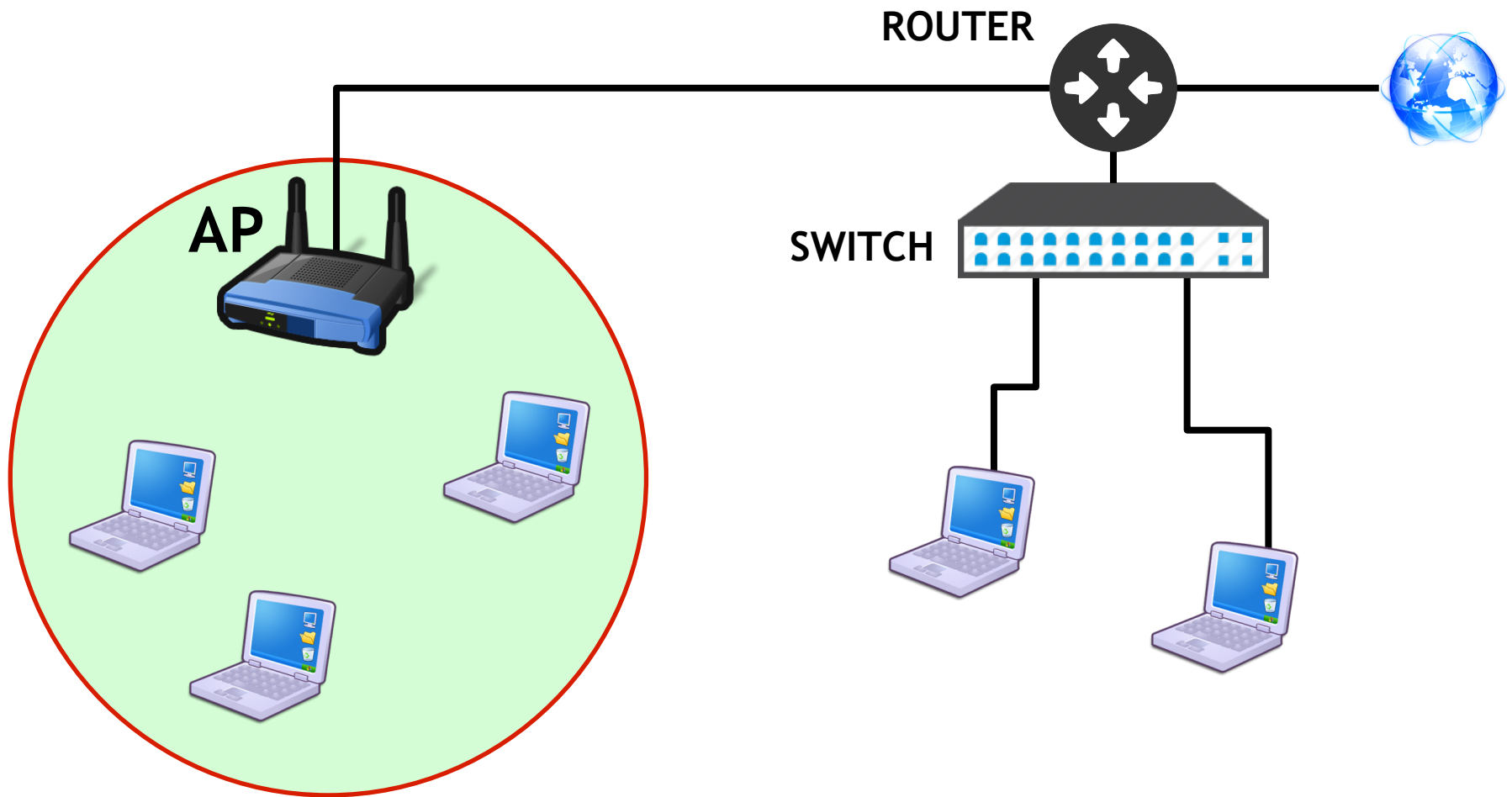
❑ Very long standard, 2012 release is approx. 2800 pages!



A LAN goes WIRELESS

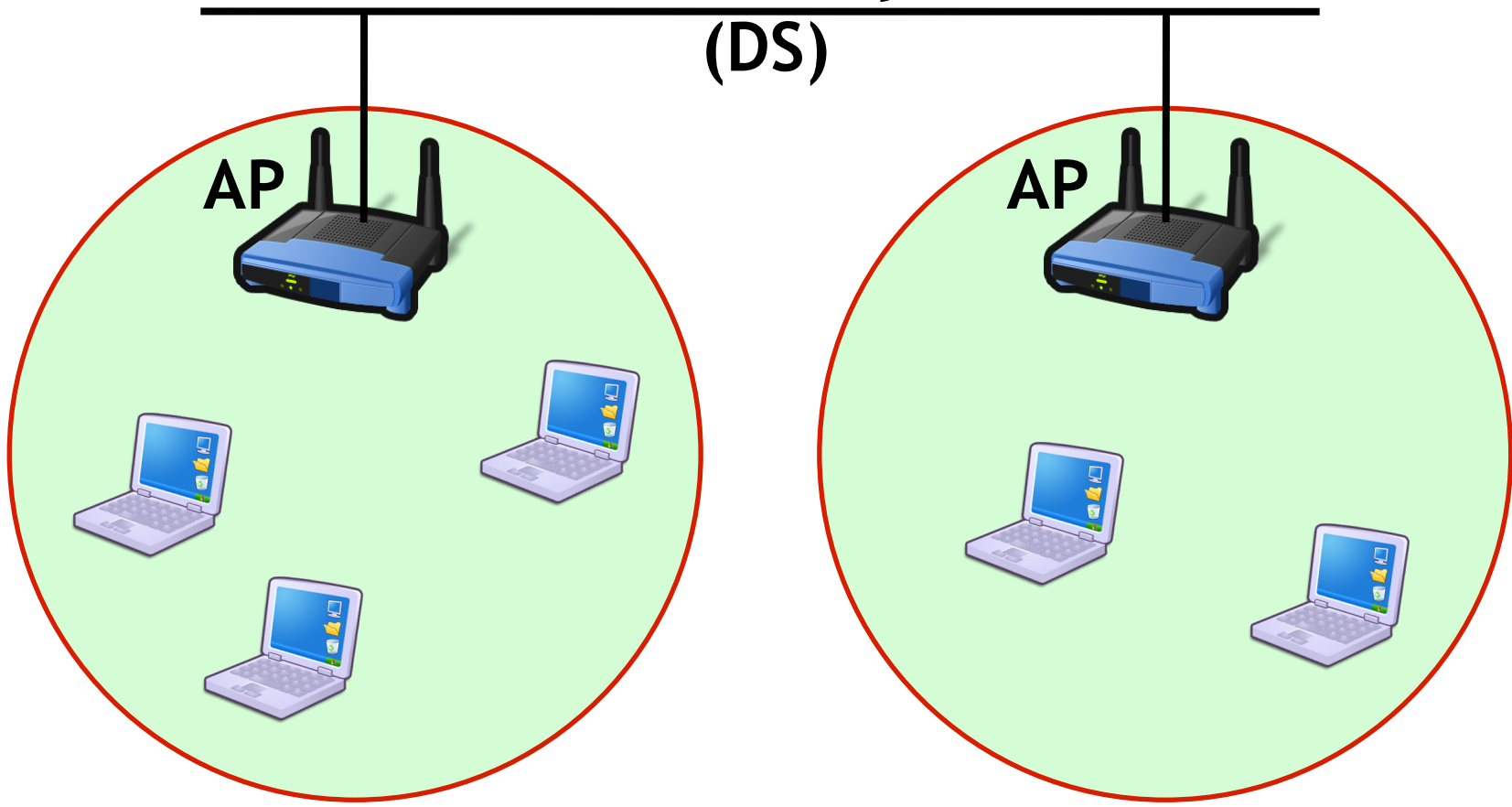


A LAN goes WIRELESS/2



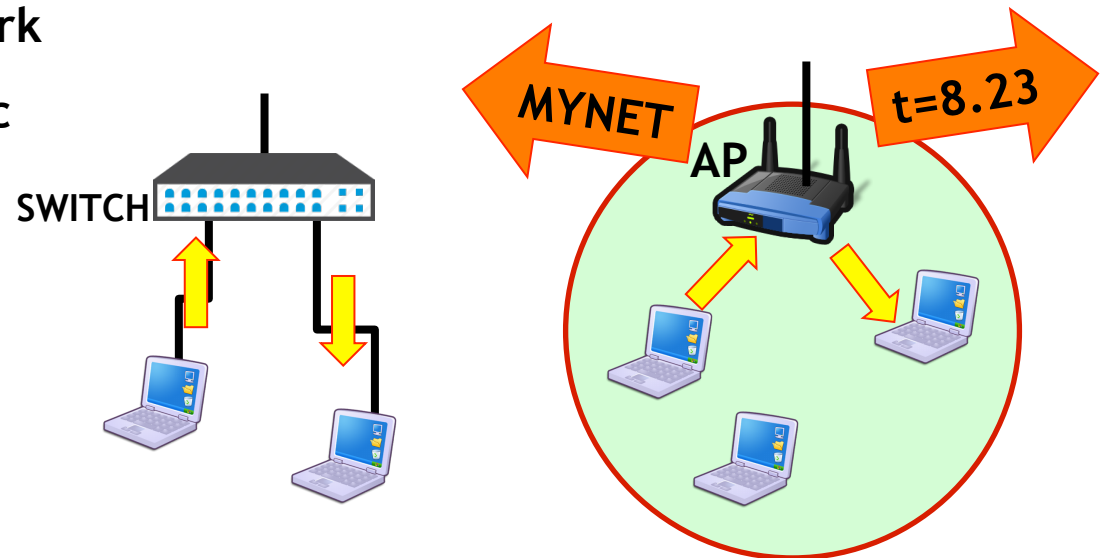
A LAN goes WIRELESS/3

Distribution System (DS)



Wireless-LAN vs Wired-Lan

- ❑ Apparently: AP replaces the switch, air replaces cables
- ❑ AP, in fact,
 - forwards inter-stations frames (no direct comm)
 - rules stations access to the network (e.g, by authenticating)
 - manages even more issues than the switch has to, i.e.,
 - advertizes the network
 - synchronizes time etc



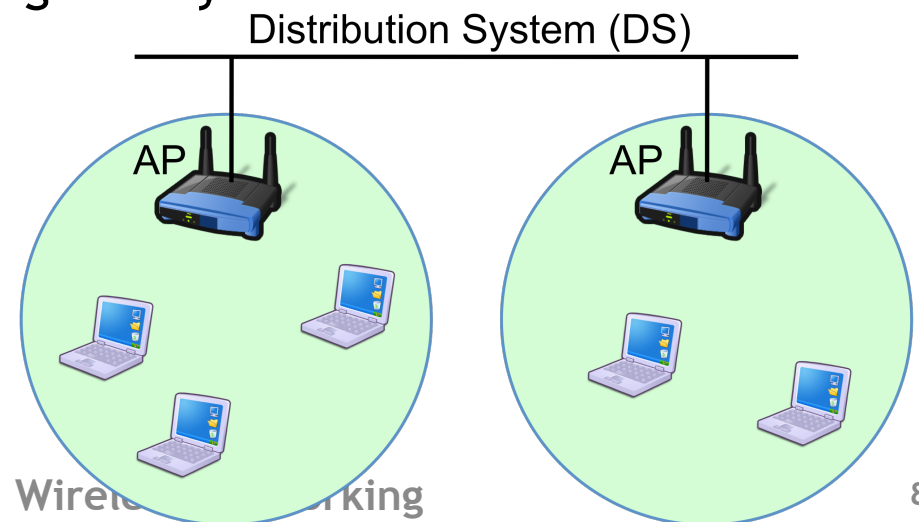
802.11 Wireless-Lan: Infrastructure Mode Basic

❑ Each cell is an Infrastructure “Basic Service Set” (BSS)

- The Access Point (AP) “creates” and maintain the BSS
- Time sync, BSS name and capabilities inside “Beacons” frame
- All BSS traffic goes through the AP

❑ More cells build an “Extended Service Set” (ESS)

- A Distribution System (DS, wired or wireless) connect all APs
- DS may connect to an Internet gateway



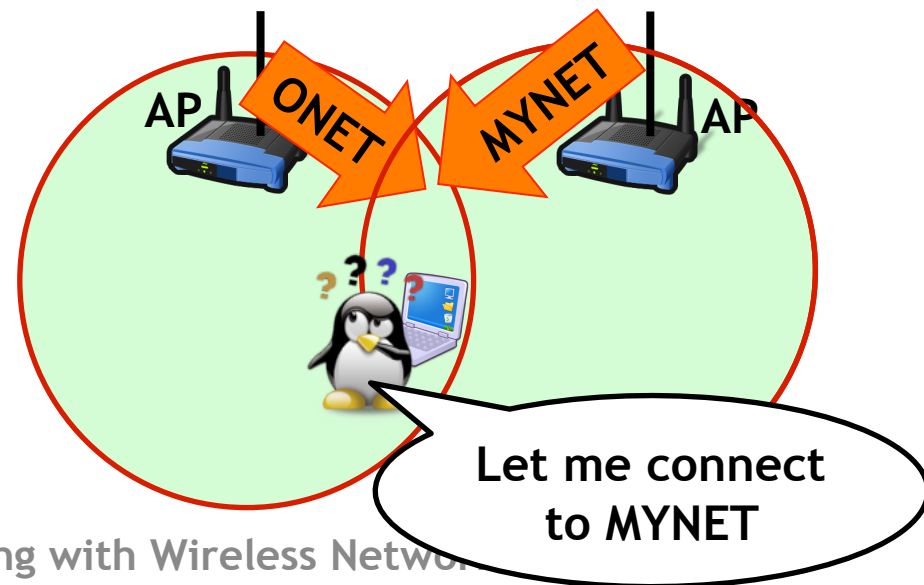
802.11 Wireless-Lan: Infrastructure Mode Basic

❑ A station that wants to connect does the following:

- Scan: check all the available networks for one known
 - Passively, by receiving “Beacons”, or Actively, by sending “Probes”
- Authenticate: proves to the AP she knows something
 - Easiest case: simply send her identity, wait for an ack
- Associate: station and AP shares mutual capabilities

❑ Eventually:

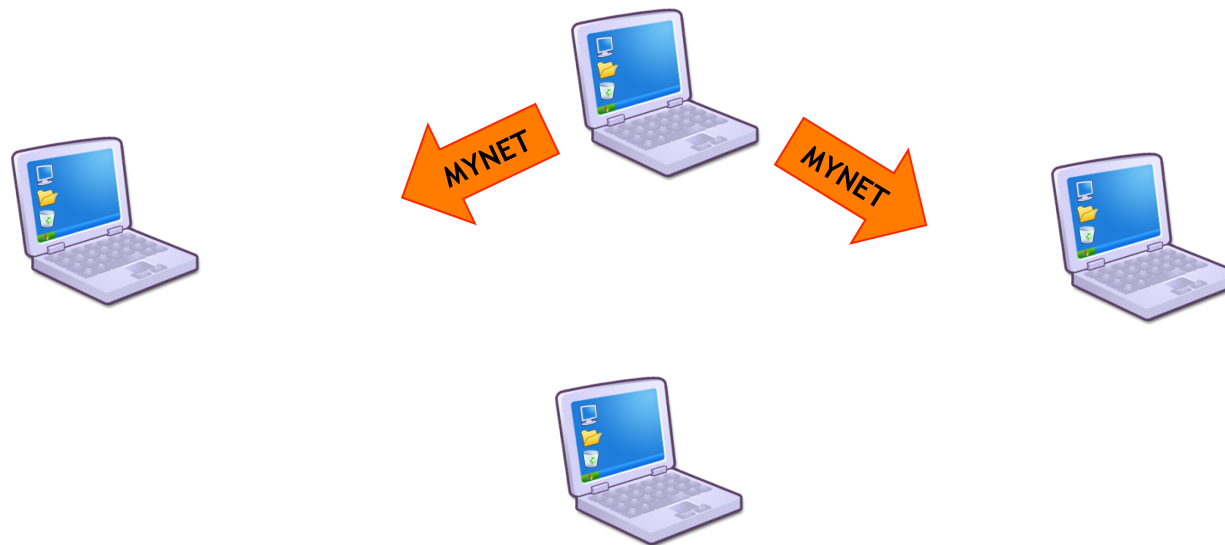
- Station sends/receives traffic



802.11 Wireless-Lan: Ad-Hoc Mode Basic

❑ Stations create an Ad-Hoc network with a specific “NAME”

- No AP needed
- The first station active starts sending beacons (e.g., leader)
- Other stations can join the Ad-Hoc network
 - If they do not receive the beacon from the leader, they transmit one!



A very incomplete standard synopsis

Document	year	modulation add-on	band (GHz)	width (MHz)	spatial stream	Rates addon
802.11	97	DSSS	2.4	20	/	1, 2
802.11b	99	CCK	2.4	20	/	5.5, 11
802.11a	01	OFDM	5	20	/	6 ... 54 (8 rates)
802.11g	03	/	2.4	20	/	all the above
802.11n	09	MIMO & OFDM+	2.4 & 5	20, 40	up to 4	HT-PHY MCS [max 600Mb/s]
802.11ac	14	MU-MIMO & OFDM++	2.4 & 5	20,40, 80,160	up to 8	VHT-PHY MCS [max 6.7Gb/s]



A very incomplete standard synopsis

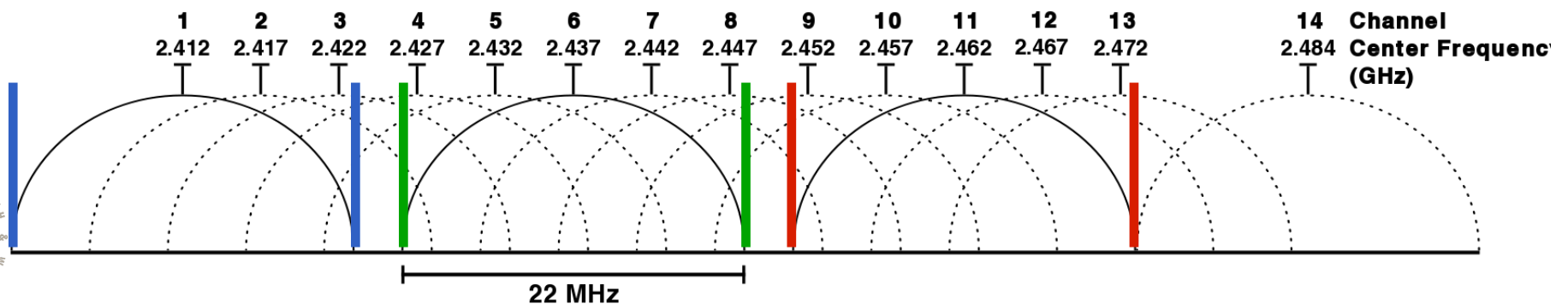
- ❑ In this course we will use 802.11bg only devices
 - Yeah, pretty old hardware but...
 - ... we have access to the NIC internals, we can play with the standard

Document	modulations	band (GHz)	width (MHz)	Rates
802.11b/g	DSSS/CCK OFDM	2.4	20	1, 2 / 5.5, 11 6, 9, 12, 18, 24, 36, 48 ,54



IEEE 802.11: insight of the 24GHz band

- ❑ ISM 2.4GHz band spans range [2400-2483.5]MHz worldwide
 - Availability subject to country regulations
 - E.g., USA [1-11], Italy [1-13], Japan none of them!
 - To make Wi-Fi working, Japan regulator allows outsider channel
- ❑ Standard: 13 channels (5 MHz spacing) + channel 14
 - ch_N @ [2407 + 5 * N]MHz, $1 \leq N \leq 13$; very busy ☹
 - ch_{14} @ 2484MHz; not used outside Japan ☺
- ❑ How many orthogonal channels? Remember 20MHz width





Wireless LAN Standard

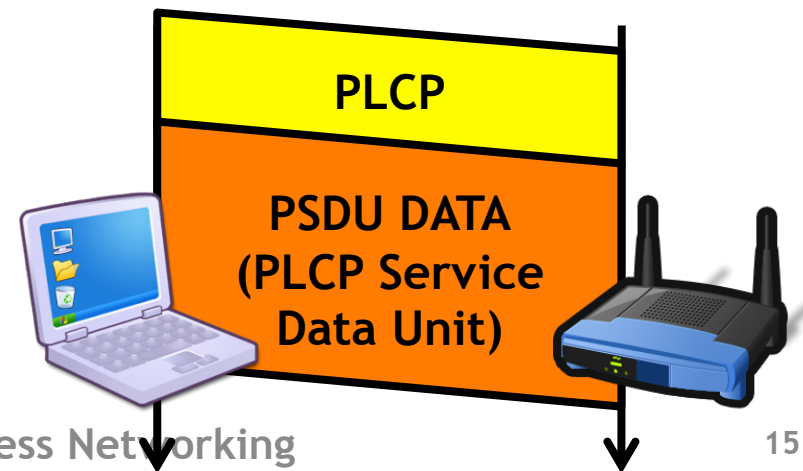
802.11bg Physical Layer analysis



IEEE 802.11bg: a frame unit (no MAC yet)

- ❑ Each frame preceded by PLCP preamble
 - Physical Layer Convergence Procedure
- ❑ PLCP helps the receiver
 - Understanding a transmission is beginning (energy raise)
 - Knowing which data-rate encoding is used for data and its length
 - Synchronizing the decoding subsystem
- ❑ PLCPs of 11b and 11g differ
 - Let's check!

PPDU
(PLCP Protocol
Data Unit)



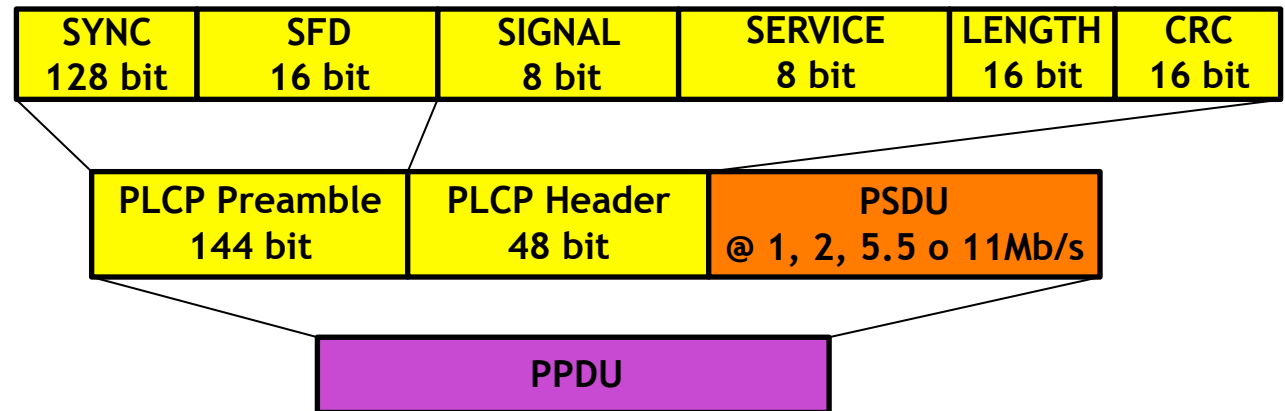
Playing with Wireless Networking



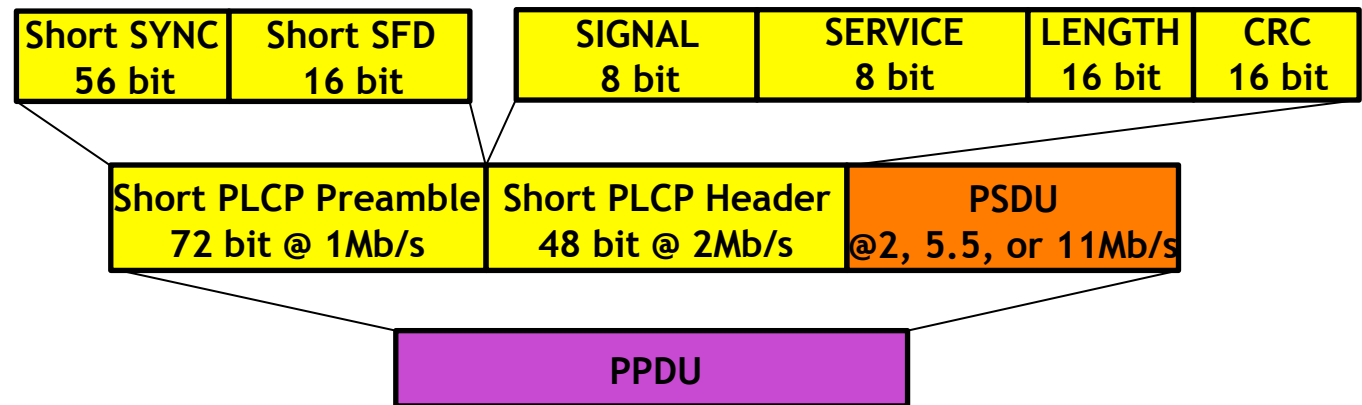
IEEE 802.11b: PPDU format

□ Two possible PLCP format:

Long PLCP is 192 μ s



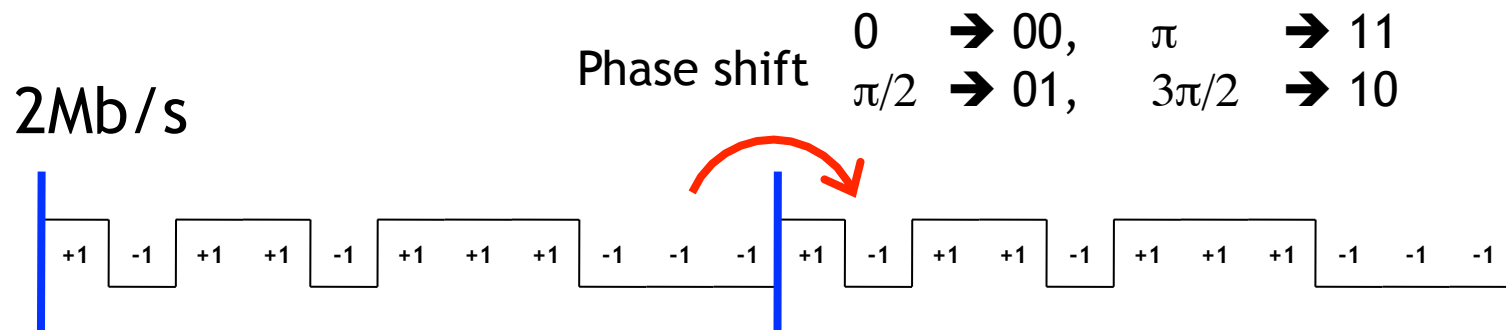
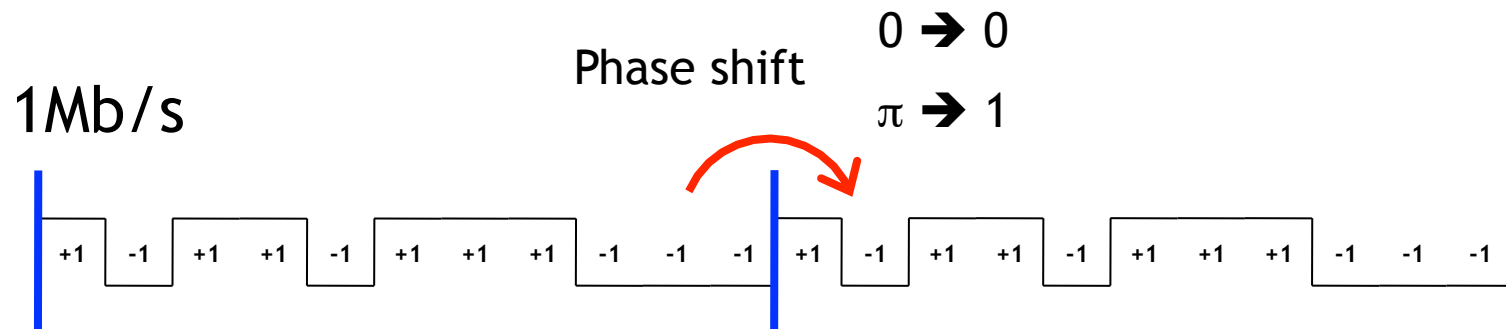
Short PLCP is 96 μ s
For 2, 5.5, 11Mb/s



IEEE 802.11b: PPDU format/2

❑ Direct Sequence Spread Spectrum modulations

- A sequence of 11 “chips” transmitted repeatedly by shifting phase
- Phase shift of consecutive chip trains encode the PPDU



IEEE 802.11b: PPDU format/3

❑ PPDU format is very inefficient!

❑ E.g.: Acknowledgement (shortest frame), 14byte = 112bit

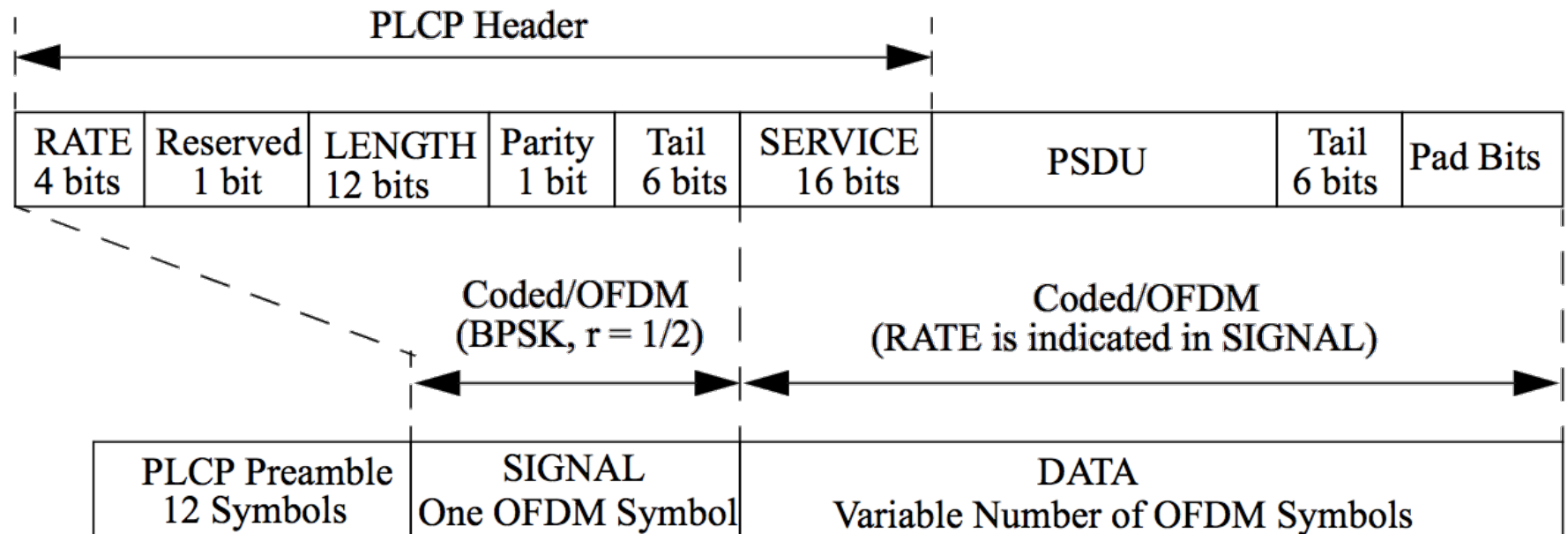
- @1Mb/s, PSDU:= $112\mu s$
- @2Mb/s, PSDU:= $56\mu s$
- @5.5Mb/s, PSDU:= $21\mu s$
- @11Mb/s, PSDU:= $11\mu s$

PLCP much longer than actual data!



IEEE 802.11g: PPDU format

- 802.11g: new PLCP, very short



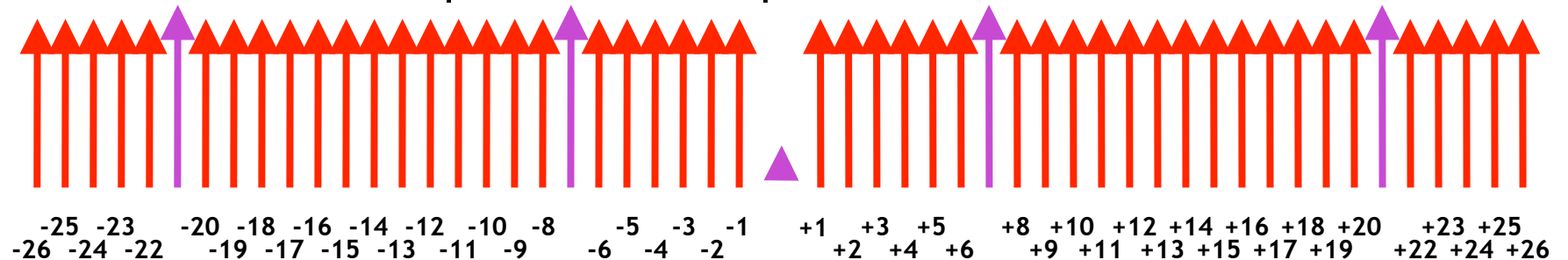
PLCP is 20 μ s



IEEE 802.11g: PPDU format/2

❑ Orthogonal Frequency Division Multiplexing (6Mb/s)

- PPDU is divided in groups of 24 bits (three bytes)
- Each group of 24 is expanded to 48 bit (FEC)
- Each bit of these 48 weights a specific carrier (with -1 for 0, +1 for 1)
- Five additional pilot carriers - pink - inserted



- Time signal computed by running ifft of the carrier weights
- Each group takes $4\mu\text{s}$, call this “OFDM symbol”

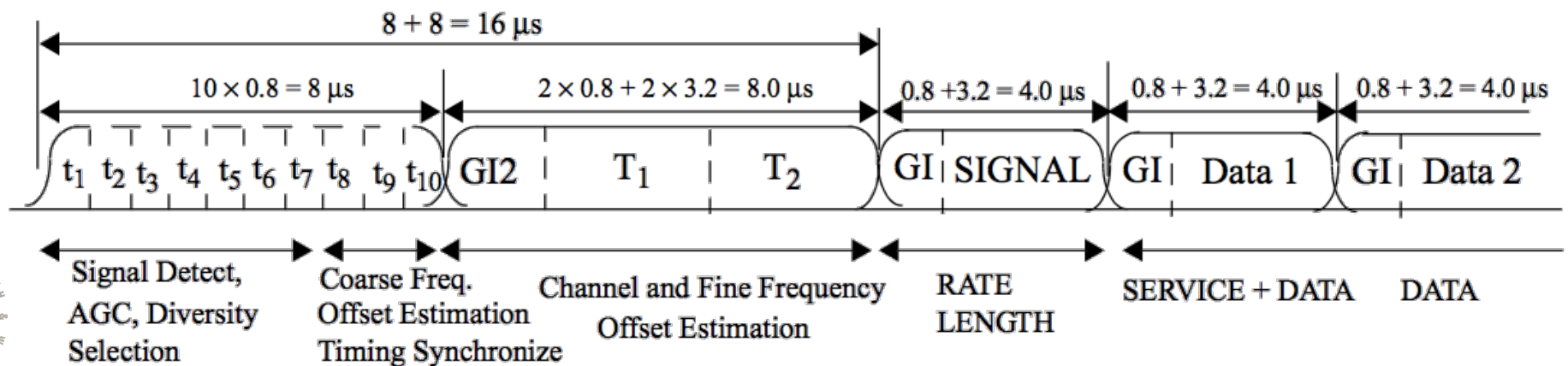
❑ Higher data rates map more bits to each carrier/symbol



IEEE 802.11g: PPDU format/2

□ 802.11g: frame includes

- PLCP Preamble made of 10 short symbols ($8\mu s$) + 2 long ones ($8\mu s$)
- PLCP Header made of
 - SIGNAL field in its own symbol ($4\mu s$, PSDU length+rate)
 - SERVICE field, first 16 bits of first data symbol
- PSDU made of symbols, each one carrying N bits, N depends on Rate
- PSDU terminate with CRC32 and at least 6 padding bits



IEEE 802.11g: PPDU format/2

□ 802.11g: data payload

- Bit expanded by convolutional encoder for FEC, $R = [1/2, 2/3, 3/4]$
- Groups of N_{CBPS} (Coded Bit Per Symbol) or N_{DBPS} (Data Bit Per Symbol)
- Each subcarrier transport N_{BPSC} bit (Bit Per SubCarrier)

Modulation	R	N_{BPSC}	N_{CBPS}	N_{DBPS}	Data rate
BPSK	1/2	1	48	24	6
BPSK	3/4	1	48	36	9
QPSK	1/2	2	96	48	12
QPSK	3/4	2	96	72	18
16-QAM	1/2	4	192	96	24
16-QAM	3/4	4	192	144	36
64-QAM	2/3	6	288	192	48
64-QAM	3/4	6	288	216	54



IEEE 802.11g: frame format/6

- ❑ E.g.: Acknowledgement, 14byte = 112bit
- ❑ PLCP: $20\mu s$
- ❑ $DATA_{PSDU}: 16b(\text{SERVICE})+112b(\text{PSDU})+6b(\text{tail}_{\min})=134b$

Data rate	PLCP	N_{DBPS}	bit	symbol	ΔT	Extension	Total
6	$20\mu s$	24	134	6	$24\mu s$	$6\mu s$	$50\mu s$
9	$20\mu s$	36	134	4	$16\mu s$	$6\mu s$	$42\mu s$
12	$20\mu s$	48	134	3	$12\mu s$	$6\mu s$	$38\mu s$
18	$20\mu s$	72	134	2	$8\mu s$	$6\mu s$	$34\mu s$
24	$20\mu s$	96	134	2	$8\mu s$	$6\mu s$	$34\mu s$
36	$20\mu s$	144	134	1	$4\mu s$	$6\mu s$	$30\mu s$
48	$20\mu s$	192	134	1	$4\mu s$	$6\mu s$	$30\mu s$
54	$20\mu s$	216	134	1	$4\mu s$	$6\mu s$	$30\mu s$





Wireless LAN Standard

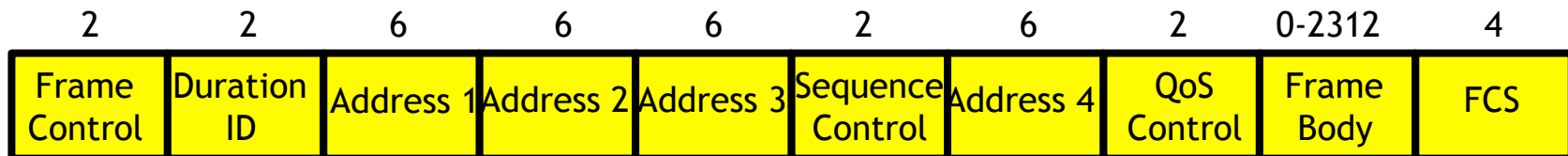
802.11bg Frame format analysis (@ layer 2)



IEEE 802.11: Frame types

□ Three types of Mac Protocol Data Unit (MPDU):

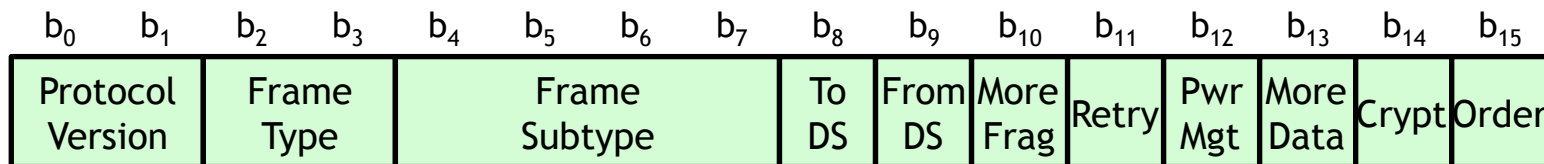
- Management, e.g. Association Request/Response, Beacon, (De)Auth
 - Network/BSS Advertisement, BSS Join, Authentication etc
- Control, e.g. ACK, RTS, CTS, Poll, etc
 - For channel access (RTS, CTS), positive frame acknowledgment
- Data: Plain data + QoS Data, etc
 - Frames with user data
- MPDU fields: depend on frame type!



IEEE 802.11: PSDU fields/1

□ Frame Control:

- Protocol version, only 0 today
- Type and Subtype encode frame type + subtype
- ToDS: frame is for Distribution System; FromDS frame is from DS
 - If both set to 1, frame is transported by a Wireless DS
- More: announce other fragments are coming (PSDU is fragmented)
- Retry: help rx'er understanding this is a retransmission
- {Pwr Mgt, More Data} deal with power management, save
- Protected: announce Frame Body is encrypted



IEEE 802.11: PSDU fields/2

❑ Duration/ID

- Meaning depends on MPDU type
- Data: number of μs after frame end during which medium is reserved
 - Used by Virtual Carrier Sense

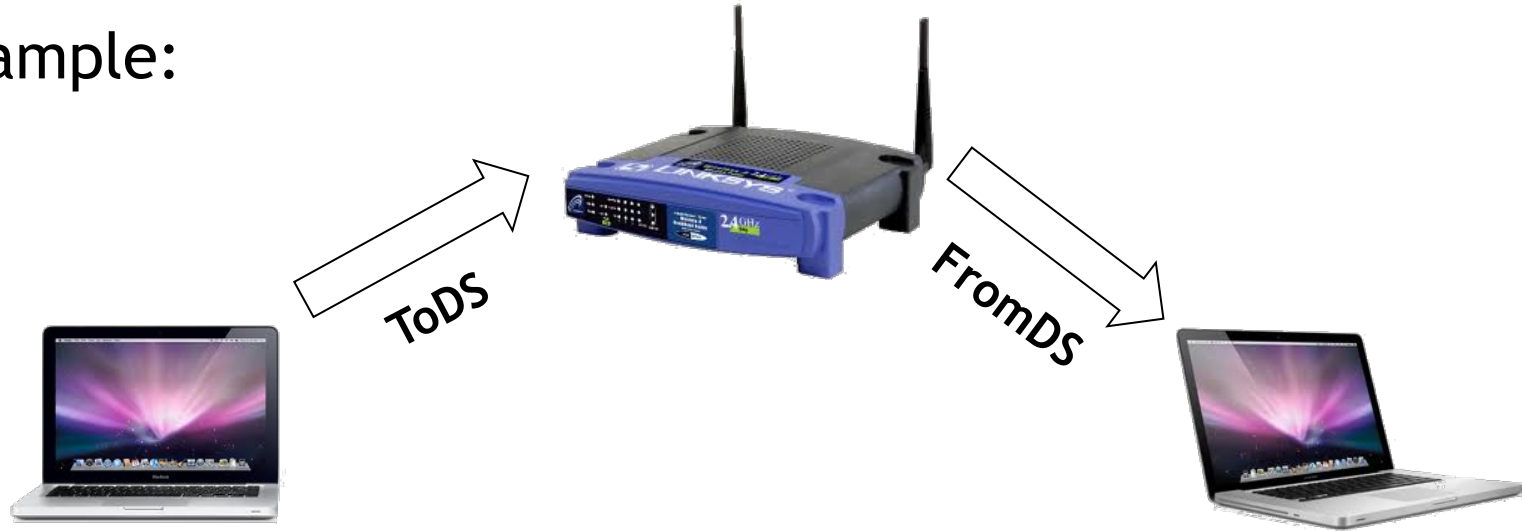
❑ Address fields: they depends on ToDS/FromDS fields:

- BSSID: Basic Service Set IDentification
 - Address of the AP
- DA: Destination Address, “final destination”
- RA: Receiver Address, immediate frame destination
- SA: Source Address, who has generated this frame
- TA: Transmitter Address, who has forwarded this frame



IEEE 802.11: PSDU fields/3

□ Example:



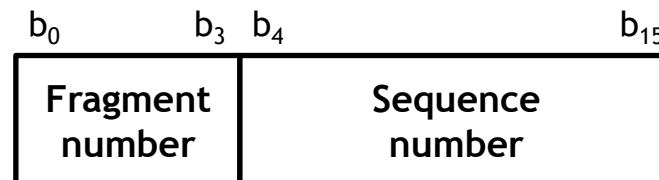
To DS	From DS	Address 1	Address 2	Address 3	Address 4
0	0	RA = DA	TA = SA	BSSID	N/A
0	1	RA = DA	TA = BSSID	SA	N/A
1	0	RA = BSSID	TA = SA	DA	N/A
1	1	RA	TA	DA	SA



IEEE 802.11: PSDU fields/4

❑ Sequence Control:

- Fragment number, 4 bits
 - For fragmented PSDU, it's the number of this fragment
- Sequence Number, 12 bits, unique for PSDU
 - Identify the PSDU (used by rx'er to avoid accepting same frame > once)



❑ QoS Control: identify Traffic Category (optional field)

❑ FCS: CRC/32 Frame Check Sequence protecting the PSDU

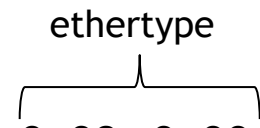


IEEE 802.11: PSDU example - Data Frame

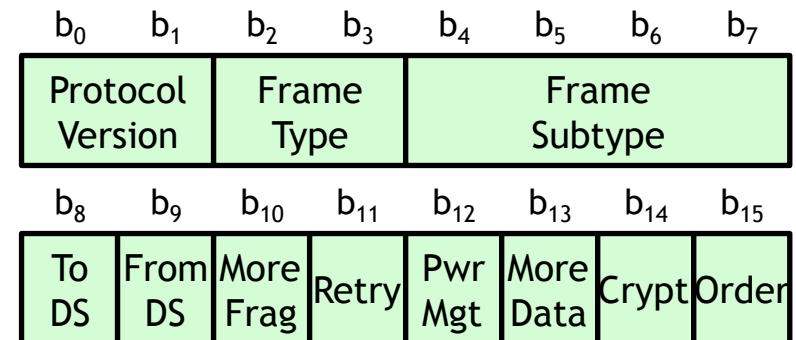
❑ IP packet, no QoS, from STA to AP (ToDS): Data frame

- Logical Link Control (LLC) encapsulation is used

- 8 bytes before IP: 0xAA, 0xAA, 0x03, 0x00, 0x00, 0x00, 0x08, 0x00



- Type: Data frame
SubType: 0 ⇒ Byte#0 := 0x08

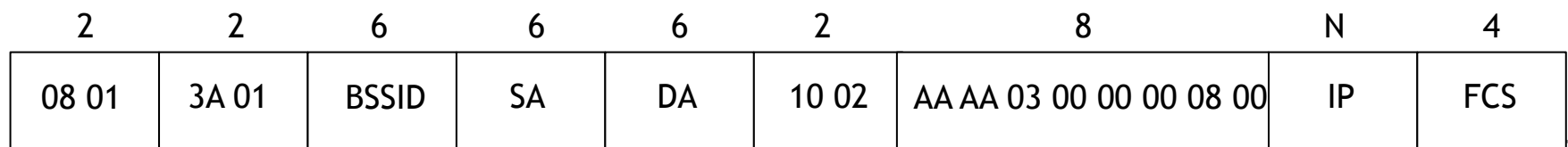
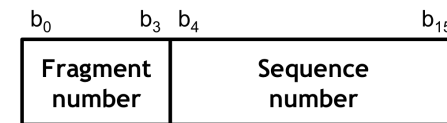


- ToDS ⇒ Byte#1 := 0x01

- Duration: time to tx an ACK + SIFS

- Address: it's a ToDS frame, fill the three address fields

- SeqCTRL: seq. no:=33 ⇒ SeqCTRL:=0x0210



IEEE 802.11: PSDU example - Data Frame

❑ Example with WireShark

- Open a trace file
- Show Beacons
- Show Data, retry etc





Wireless LAN Standard

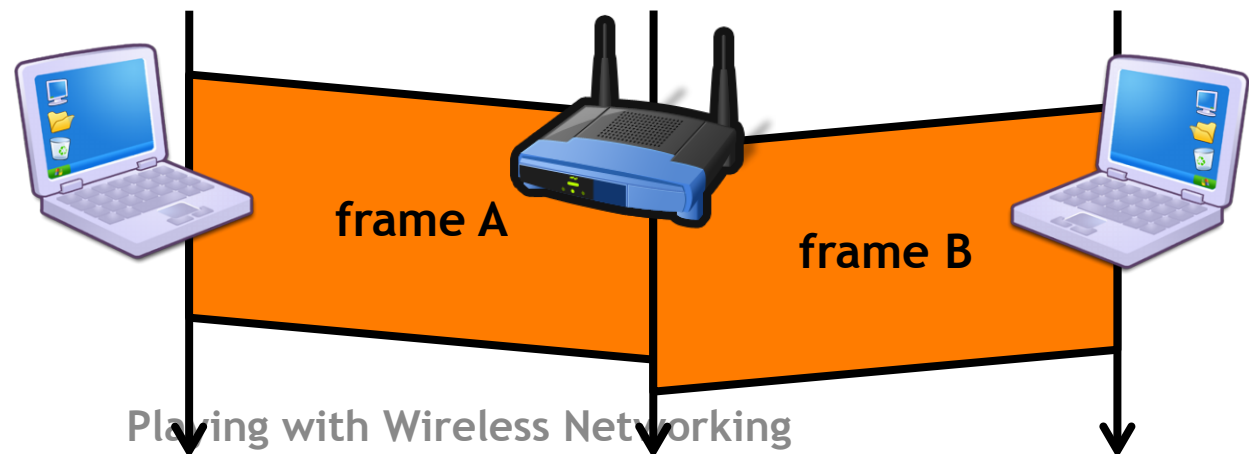
The Basic Access Scheme

Distributed Coordination Function (DCF)



Transmissions in a Broadcast Medium

- ❑ If two stations transmit at the same time
 - Receiver(s) cannot decode packets (collision)
 - Transmitters do not know whether data was received
- ❑ 802.11 standard introduces
 - Slotted medium
 - Positive Acknowledgment with Retransmission
 - Carrier Sense Multiple Access/Collision Avoidance



Time Slot

- ❑ Time is divided into intervals, called **slots**
 - Working with slot (tx may start with slot) reduces uncertainty

- ❑ A Slot is the system unit time
 - 802.11b Slot Time is $20\mu\text{s}$, 11g is $9\mu\text{s}$

- ❑ Time synchronized with Beacons transmitted by the BSS AP
 - Each Beacon carries a 64bit time value ($1\mu\text{s}$ granularity)
 - Stations in the BSS copy beacon time to their clock registers
 - Skews due to poor clock design periodically corrected

- ❑ A BSS is a synchronous system!!



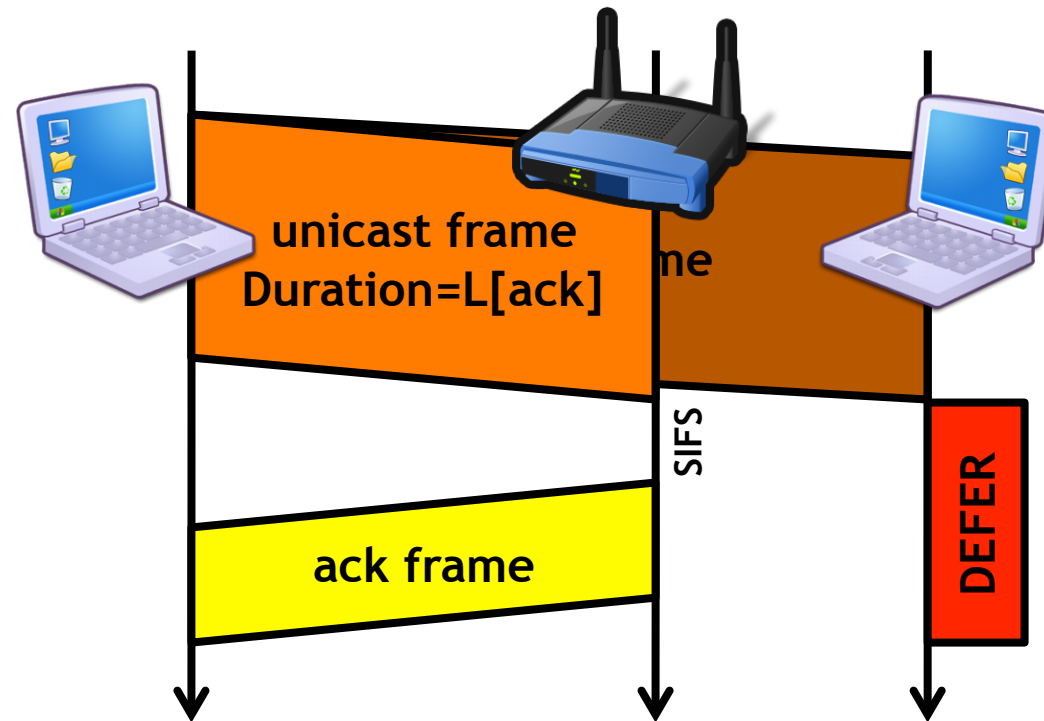
InterFrame Space (IFS)

- ❑ Time interval between consecutive transmissions
- ❑ Different IFSs allow different access priorities
 - Short IFS: separate transmissions belonging to the same dialogue
 - SIFS in 11bg it's $10\mu\text{s}$
 - Point Coordination IFS: used by the Point Coordinator
 - PIFS is SIFS + Slot Time
 - Distributed IFS: waited by stations when contending for a free channel
 - DIFS is SIFS + $2 * \text{Slot Time}$
 - Extended IFS: waited by stations when receiving a bad frame
 - EIFS is SIFS + TxTime[AckFrame] + DIFS



Positive Acknowledgment & Retransmissions

- ❑ Received unicast frame must be acknowledged
 - Other nodes defer transmissions by using the received “Duration”



Playing with Wireless Networking

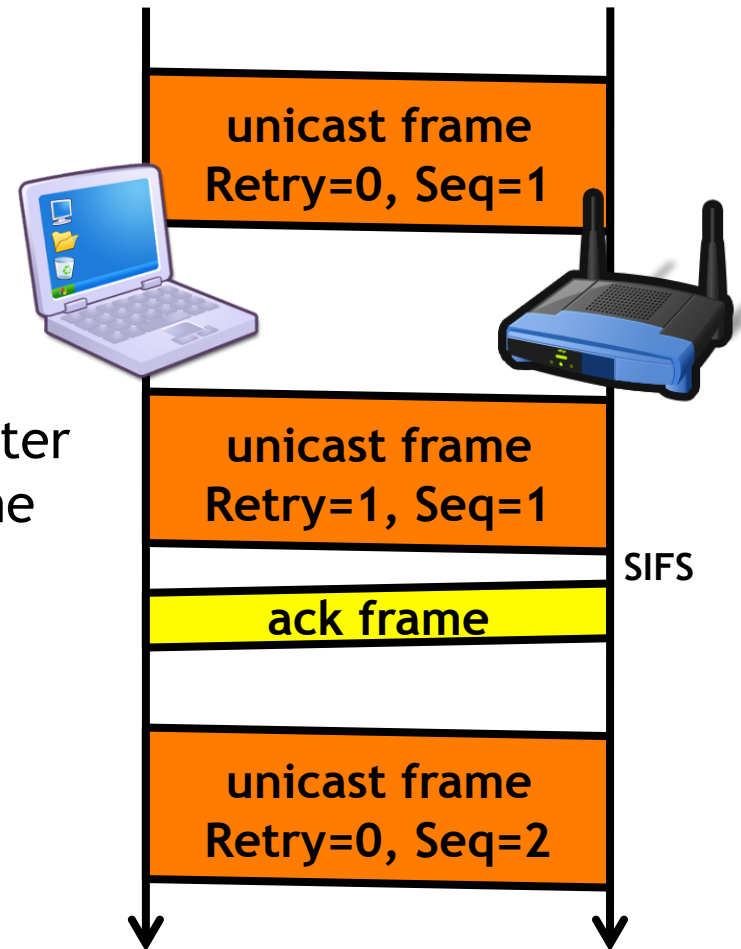
Positive Acknowledgment & Retransmissions/2

❑ If no acknowledgment coming from receiver

- Retransmit the frame with Retry bit set and same sequence counter

❑ When acknowledgment received

- Increase the sequence counter and switch to the next frame



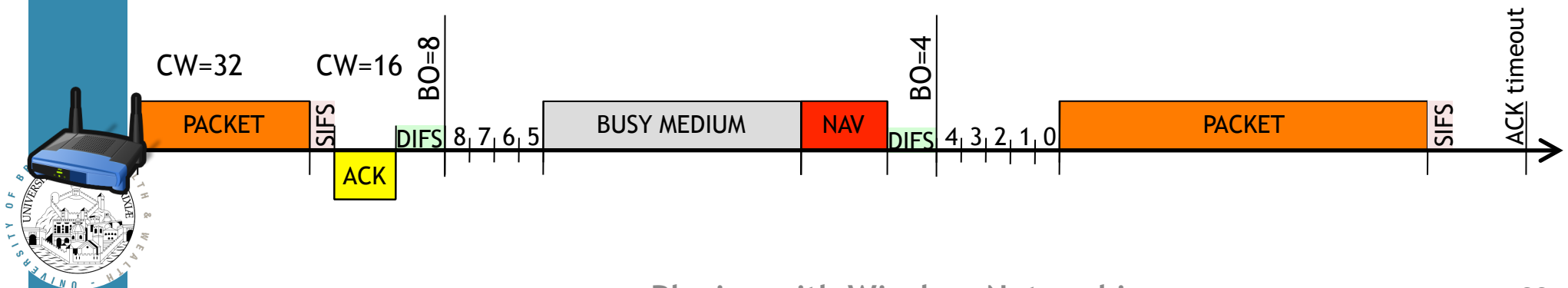
Carrier Sense Multiple Access/ Collision Avoidance

- ❑ Apart from slot synchronization
 - No other explicit coordination among stations
- ❑ To **avoid** repeated collisions
 - Carrier Sense Multiple Access/Collision Avoidance (CSMA/CA)
 - Given free-space attenuation, /Collision Detection is not feasible!
- ❑ CSMA/CA Basic
 - Stations willing to transmit have to contend for channel access
 - A station repeats the contention procedure for every (re)transmission

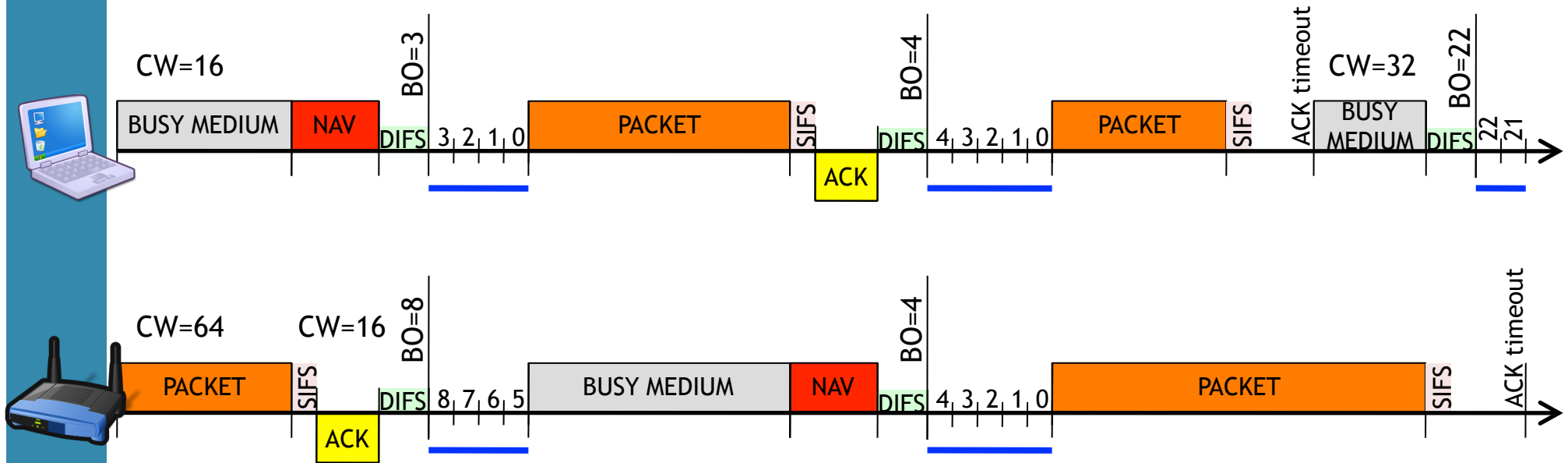


CSMA/CA and channel contention

- ❑ Each station keeps a Contention Window (CW) parameter
 1. At the end of the previous transmission attempt
 - If collision (no ack), double CW, otherwise reset to CW_{min}
 - Extract Backoff value $(BO) \in U[0, CW - 1]$
 2. “Monitor channel free for $t > DIFS$ ”
 3. Backoff stage: decrement BO to zero
 - Backoff: if medium free, decrement BO at every SLOT
 - When medium busy \Rightarrow Suspend: BO freezed & goto 2
 - If $BCKOFF == 0 \Rightarrow$ Transmit!



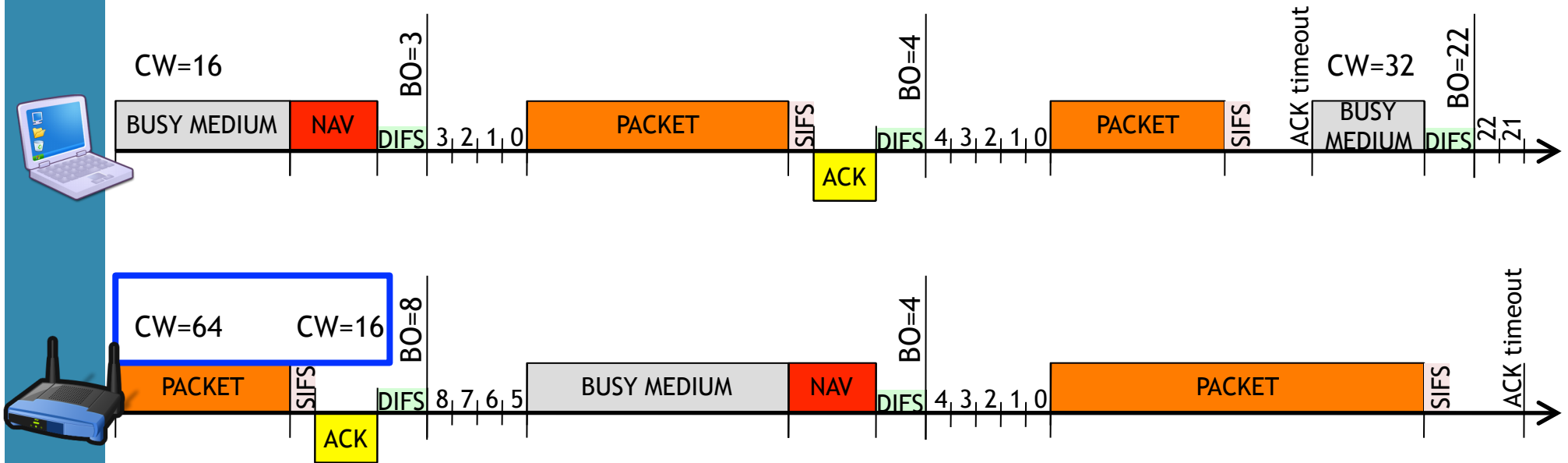
CSMA/CA and channel contention/2



Backoff count-down, channel idle



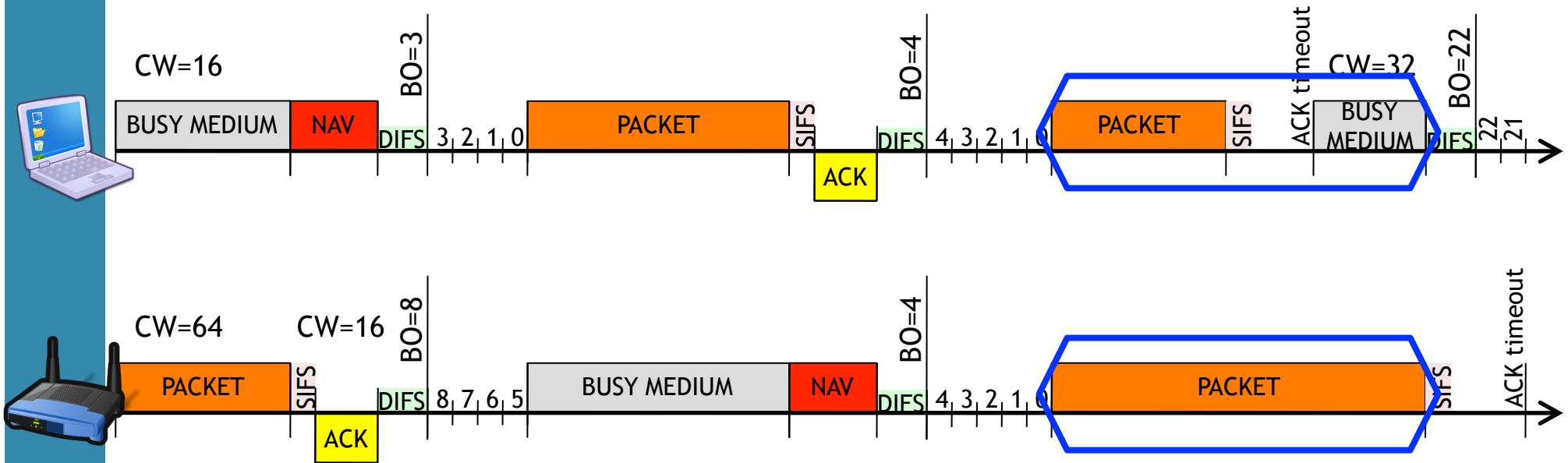
CSMA/CA and channel contention/3



CW reset after success



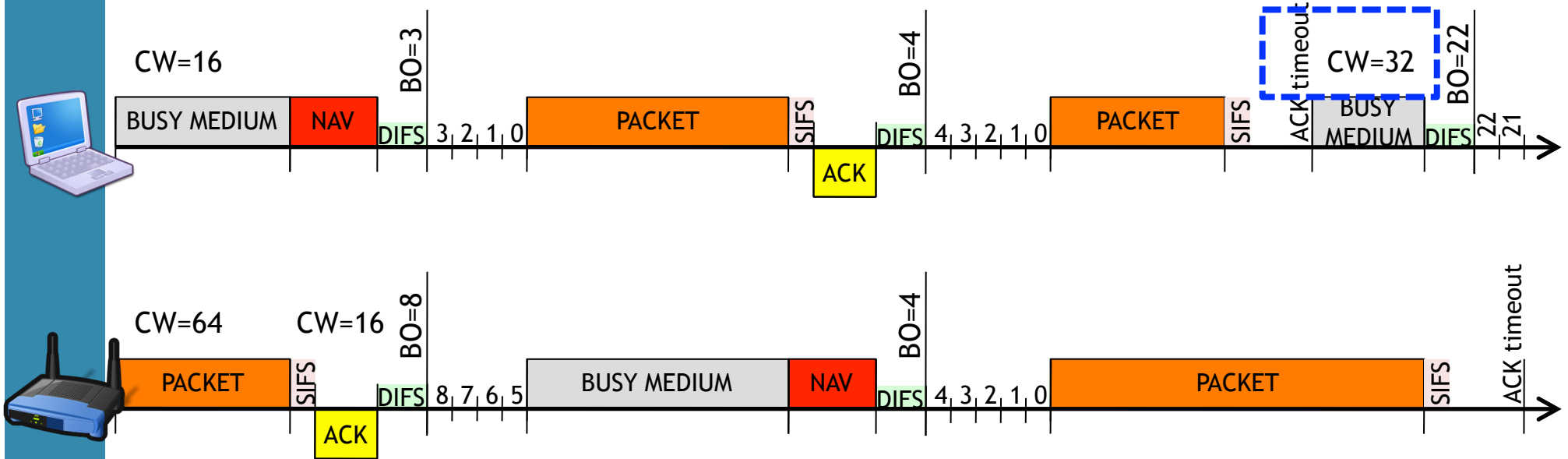
CSMA/CA and channel contention/4



Collision!



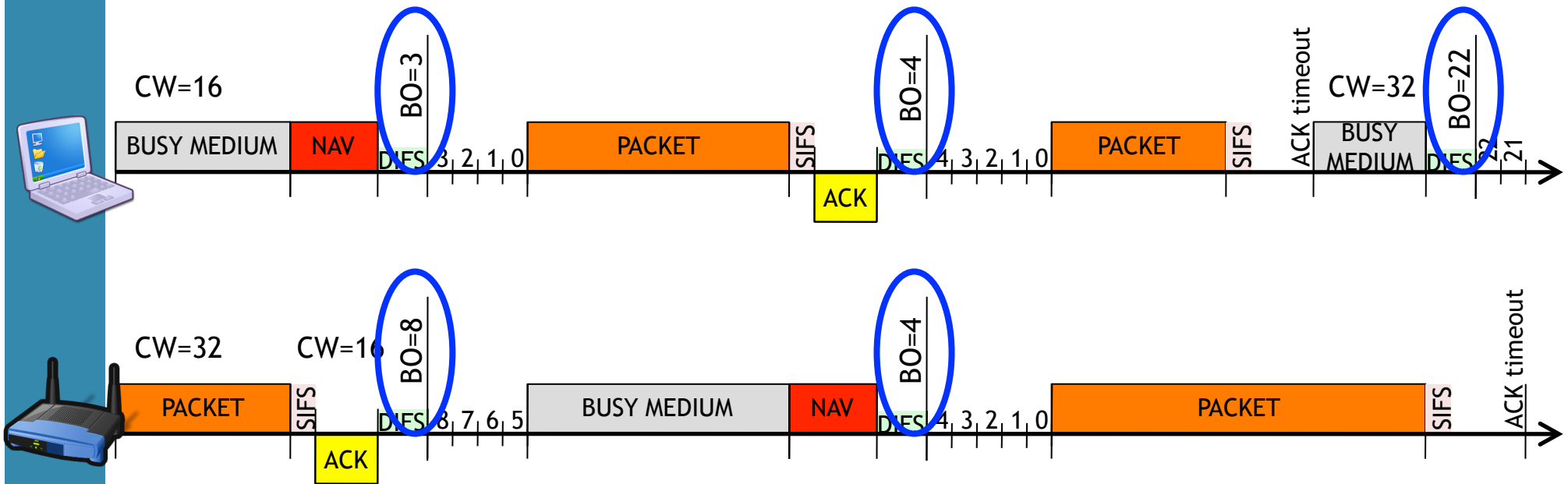
CSMA/CA and channel contention/5



CW doubling after failure



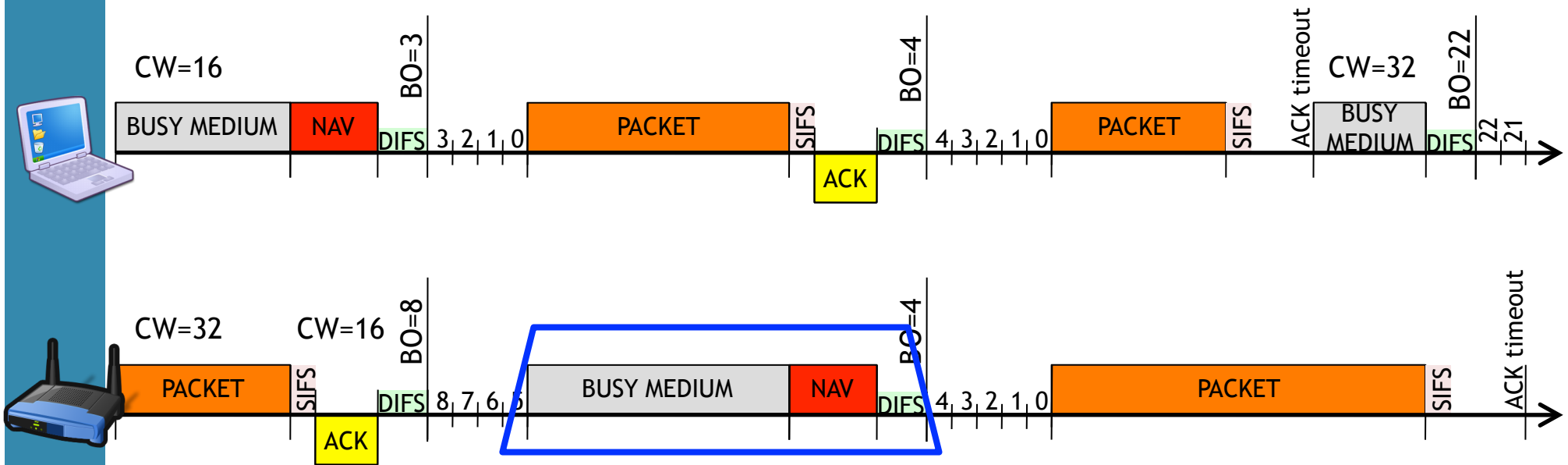
CSMA/CA and channel contention/6



Backoff Extraction



CSMA/CA and channel contention/7



**Backoff Freeze,
channel busy**



CSMA/CA, Exponential Backoff rule

- ❑ BCKOFF value is computed after every tx attempt
 - BCKOFF taken from $[0, 1, \dots, CW-1]$ with uniform distribution
- ❑ Contention Window (CW) refreshed
 - $CW = 2 * CW$ if after tx attempt there is a collision
 - Up to CW_{max} , then stay with CW_{max}
 - $CW = CW_{min}$ if after tx attempt, tx was acked by acknowledgment
- ❑ Standard values:
 - $CW_{min} = 16/32$, $CW_{max} = 1024$
- ❑ For tx a packet that requires ACK
 - Repeat access procedure up to MAX_{times} (e.g., 7), then discard packet
- ❑ This procedure guarantees network works correctly!!



CSMA/CA, pseudo-code

❑ Neglecting initializations:

procedura di trasmissione

```
while true do
  wait for packet
  wait for channel_idle_for_DIFS
  while channel_is_idle do ←
    BCKOFF--
    if BCKOFF == 0 then
      send packet
      wait for acknowledgement or timeout
      if ack received then
        reset CW
        remove packet
      else
        grow CW
      end if
      BCKOFF = extract CW
    leave do
  end fi
end do
end do
```





Wireless LAN Standard

Rate control algorithm (super-quick)



IEEE 802.11bg: rate choice

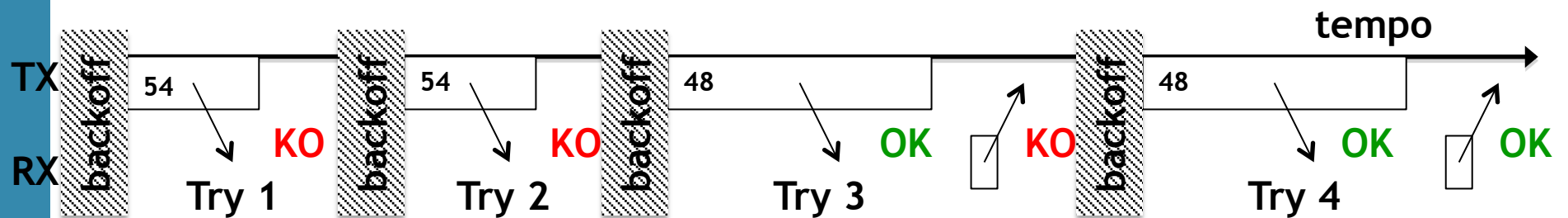
- ❑ How to choose the rate is not specified by the standard
 - Rate Controller algorithm: RC
- ❑ RCs use feedback based techniques
- ❑ E.g. Minstrel algorithm, the default today in Linux kernel
 - Count total frames transmitted PER every rate, assess success probability
 - Rate that has best success delivery ratio is the winner
 - Periodically (every N frames) send a frame at a “look-around” rate
 - Constantly scan the entire rate set
 - Rely on frames that require ACK, by counting:
 - Number of attempts per packet
 - Failed rate, success rate



IEEE 802.11: rate choice/2

□ Example: UDP packet

- RC set up these rates: $[54\text{Mb/s}^{\{1,2\}}, 48\text{Mb/s}^{\{3,4\}}, 12\text{Mb/s}^{\{5\}}, 1\text{Mb/s}^{\{6,7\}}]$



- At the end of this packet, RC refreshes its table...

Rate	Success	Failure	Rate	Success	Failure
54	2812/3004 (93%)	192/3004 (7%)	54	2812/3006 (93%)	194/3006 (7%)
48	408/507 (80%)	99/507 (20%)	48	409/509 (80%)	100/509 (20%)
36	102/402 (25%)	300/402 (75%)	36	102/402 (25%)	300/402 (75%)

- Don't change decision (not now 😊)



Bibliography

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